

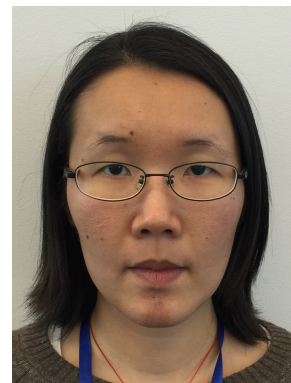
# Spectral Analysis of Stochastic Networks

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(joint work with Tingyue Gan (AMSC, UMD))



Workshop: Fluctuation—driven phenomena in non-equilibrium statistical mechanics,  
Warwick University, UK, September 2015

Networks with pairwise rates of the form  
 $L_{ij} = a_{ij} \exp(-U_{ij} / T)$ ,  $T =$  a small parameter (temperature)

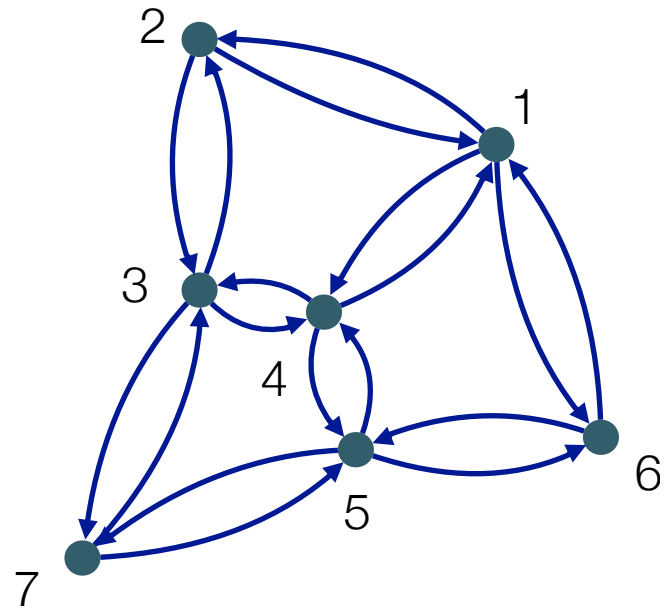
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## Time - reversible

1. Stochastic networks representing energy landscapes of atomic and molecular clusters and proteins (Wales's group (Cambridge Univ.))
2. Evolutionary genetics: fitness landscapes. (Kimura, Ewens, Gillespie) (A. Morozov & M. Manhart)
3. Markov State Model in Molecular Dynamics. (Schuette, Swope, Pande, Noe, etc)

## Time - irreversible

1. Molecular motors. (Astumian (2005))
2. Aggregation of interacting particles.
3. Stochastically oscillating energy landscapes.

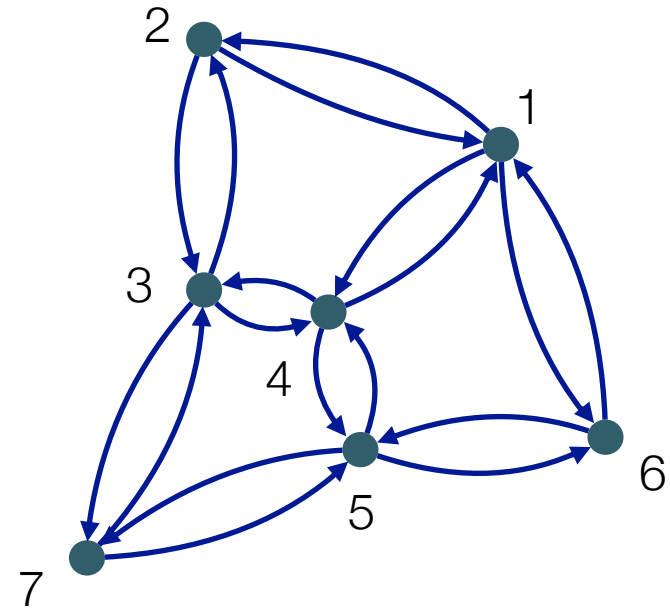
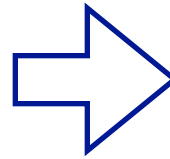
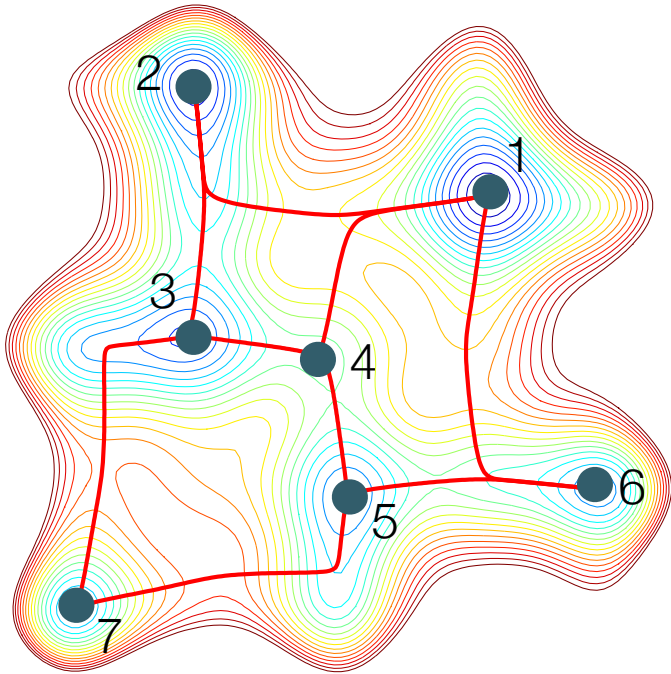


# Reversible networks

$$U_{ij} = V_{ij} - V_i,$$

$$c_{ij} = \frac{c_{ij}}{c_i},$$

where  $V_{ij} = V_{ji}$ ,  $k_{ij} = k_{ji}$



The generator matrix:

$$L_{ij} = \begin{cases} (k_{ij}/k_i)e^{-(V_{ij}-V_i)/T}, & i \sim j \\ 0, & \text{otherwise} \end{cases}$$

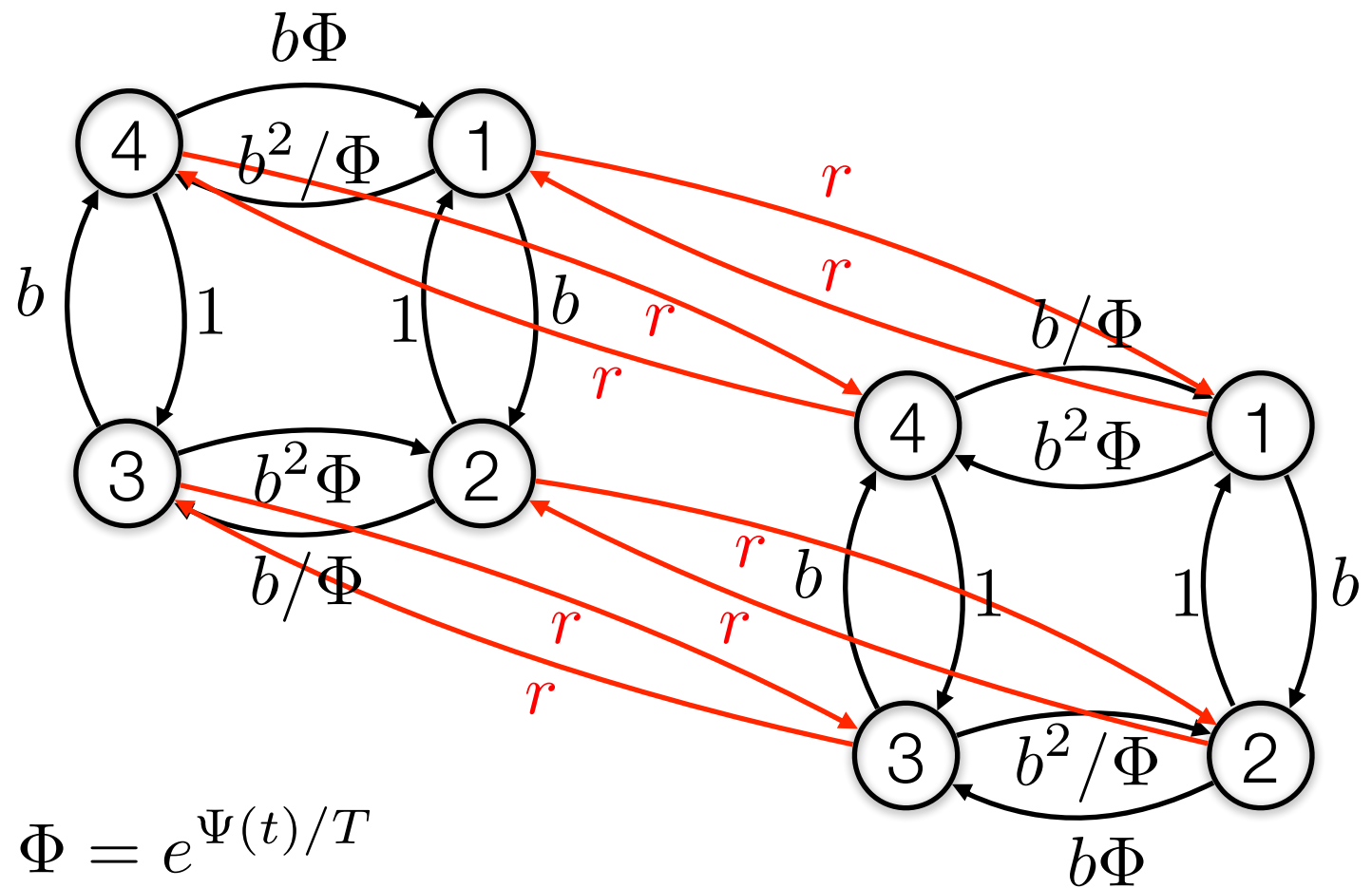
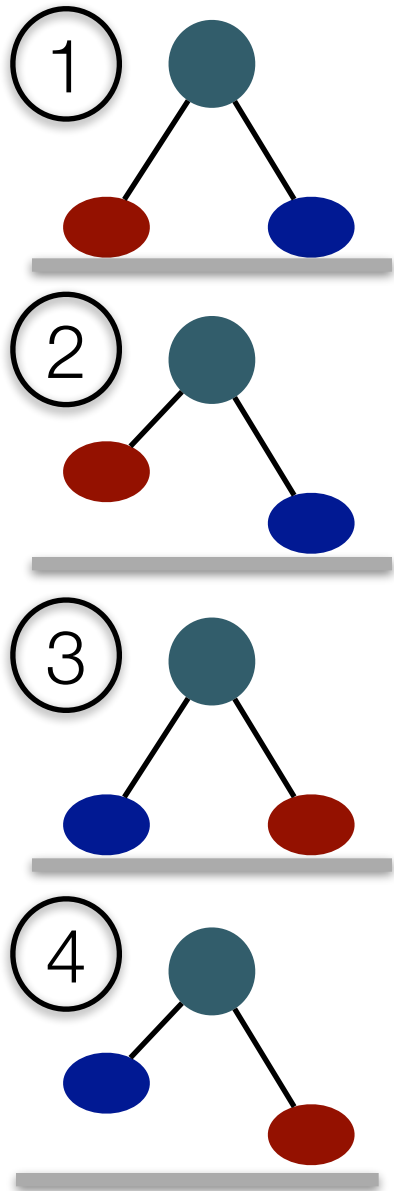
$$L_{ii} = -\sum_{j \neq i} L_{ij}$$

The invariant distribution:

$$\pi_i = k_i e^{-V_i/T}$$

# Irreversible networks: Molecular motors

(Astumian, 2005)



$$\Phi = e^{\Psi(t)/T}$$

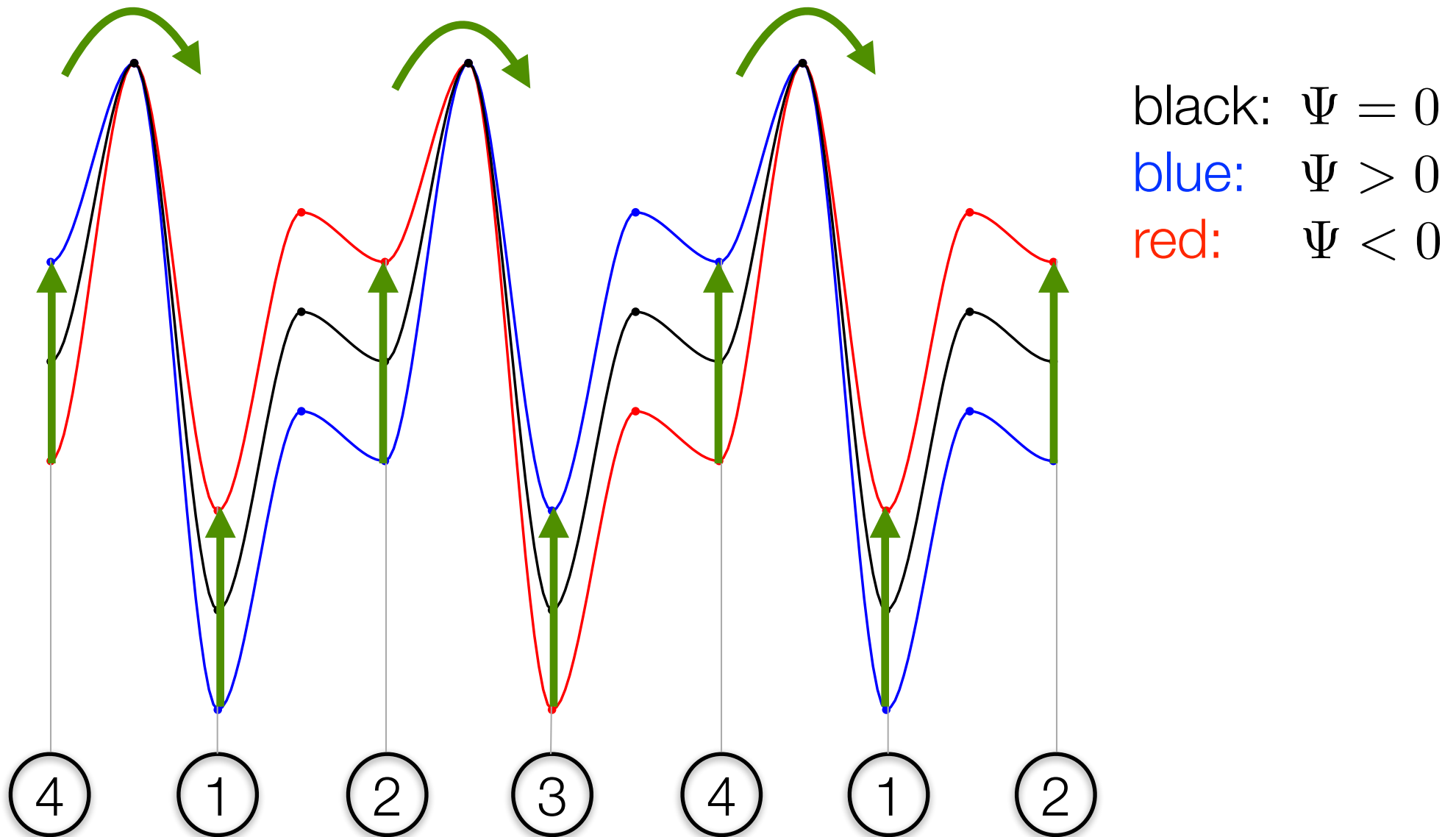
$\Psi$  switches between  $+\Psi$  and  $-\Psi$

with Poisson-distributed lifetime  $r = Ae^{-a/T}$



# Molecular motors, cont'd (Astumian, 2005)

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# Significance of the spectral decomposition

$$L = \Phi \Lambda \Psi \equiv \Phi \Lambda \Phi^{-1}$$

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- Understanding of the dynamics of the network: decompose it into a collection of processes with various time scales
- Extraction of quasi-invariant sets of states
- Building coarse-grained models

# Interpretation of spectral decomposition

$$L = \Phi \Lambda \Psi \equiv \Phi \Lambda \Phi^{-1}$$

$$\Psi = \begin{bmatrix} \pi & \rightarrow \\ \psi_1 & \rightarrow \\ \cdots & \cdots \\ \psi_{n-1} & \rightarrow \end{bmatrix} \quad \text{The matrix of left eigenvectors,}$$

$$\pi = [\pi_1, \dots, \pi_n] = \text{the invariant distribution}$$

$$\Phi = \begin{bmatrix} e & \phi_1 & \cdots & \phi_{n-1} \\ \downarrow & \downarrow & \cdots & \downarrow \end{bmatrix} \quad \text{The matrix of right eigenvectors,}$$

$$e = \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}$$

$$\Lambda = \begin{bmatrix} 0 & & & \\ & z_1 & & \\ & & \ddots & \\ & & & z_{n-1} \end{bmatrix} \quad \text{Eigenvalues: } z_k = -\lambda_k + i\mu_k$$

$$0 < \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_{n-1}$$

**The Fokker-Planck equation  
or the Master equation**

$$\frac{dp(t)}{dt} = p(t)L, \quad p(0) = p^0 = [p_1, \dots, p_n]$$

**The time evolution of  
the probability distribution**

$$p(t) = p^0 e^{tL} = p^0 \Phi e^{t\Lambda} \Psi = \pi + \sum_{k=1}^{n-1} (p_0 \phi_k) e^{z_k t} \psi_k$$

# Time evolution of the probability distribution

$$p(t) = \pi + \sum_{k=1}^{n-1} \underbrace{(p_0 \phi_k)}_{\substack{\text{Projection of} \\ \text{the initial distribution} \\ \text{onto right eigenvector}}} e^{-\lambda_k t + i\mu_k t} \underbrace{\psi_k}_{\text{Left eigenvector}}$$

Left eigenvector  $\psi_k$  decays **uniformly** across the network with rate  $\lambda_k$

## For time-reversible networks:

$L = P^{-1}Q$ , where  $P = \text{diag}\{\pi_1, \dots, \pi_n\}$ ,  $Q$  is symmetric

Right eigenvectors:  $\Phi = [\phi_0, \dots, \phi_{n-1}]$

Left eigenvectors:  $P\Phi = [P\phi_0, \dots, P\phi_{n-1}]$

Right eigenvector  $\alpha_k \phi_k$  = proportion by which the states are over- or under-populated in the perturbed distribution  $\pi + \alpha_k P\phi_k$

# Difficulties in computing spectral decomposition

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- The generator matrix is large:  $L$  in  $n$ -by- $n$ ,  $n = 10^p$ ,  $p = 4, 5, 6, \dots$
  - The pairwise rates  $U_{ij}$  can vary by tens of orders of magnitude
  - No special structure
  - Even if we succeed, the results would be hard to interpret
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## Idea

Step 1: compute the asymptotic spectral decomposition

Step 2: continue eigenpairs of interest to finite temperatures

# The goal

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- To develop an efficient algorithm for computing the zero-temperature asymptotics for eigenvalues and eigenvectors of the generator matrices
- To develop efficient continuation techniques
- Applications to large and complex networks representing energy landscapes

Asymptotics for eigenvalues and eigenvectors

# Asymptotics for eigenvalues

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## A. Wentzell, 1972

For a continuous-time Markov chain with pairwise rates of the form

$$L_{ij} \sim e^{-U_{ij}/T}$$

if all optimal W-graphs are unique, eigenvalues of the generator matrix are

$$0 > -\lambda_1 \geq \dots \geq -\lambda_{N-1}$$

$$\lambda_k \asymp \exp(-\Delta_k/T)$$

$$\Delta_k = V^{(k)} - V^{(k+1)}$$

$$V^{(k)} = \sum_{(i \rightarrow j) \in g_k^*} U_{ij}$$

where  $g_k^*$  is the optimal W-graph with  $k$  sinks

## T. Gan, C., 2015

For a continuous-time Markov chain with pairwise rates of the form

$$L_{ij} = a_{ij} e^{-U_{ij}/T}$$

if all optimal W-graphs are unique, eigenvalues of the generator matrix are

$$0 > -\lambda_1 \geq \dots \geq -\lambda_{N-1}$$

$$\lambda_k = A_k \exp(-\Delta_k/T)$$

$$\Delta_k = V^{(k)} - V^{(k+1)}$$

$$V^{(k)} = \sum_{(i \rightarrow j) \in g_k^*} U_{ij}$$

$$A_k = \frac{\prod_{(i \rightarrow j) \in g_k^*} U_{ij}}{\prod_{(i \rightarrow j) \in g_{k+1}^*} U_{ij}} + o(1)$$



# W-graphs (Wentzell, 1972)

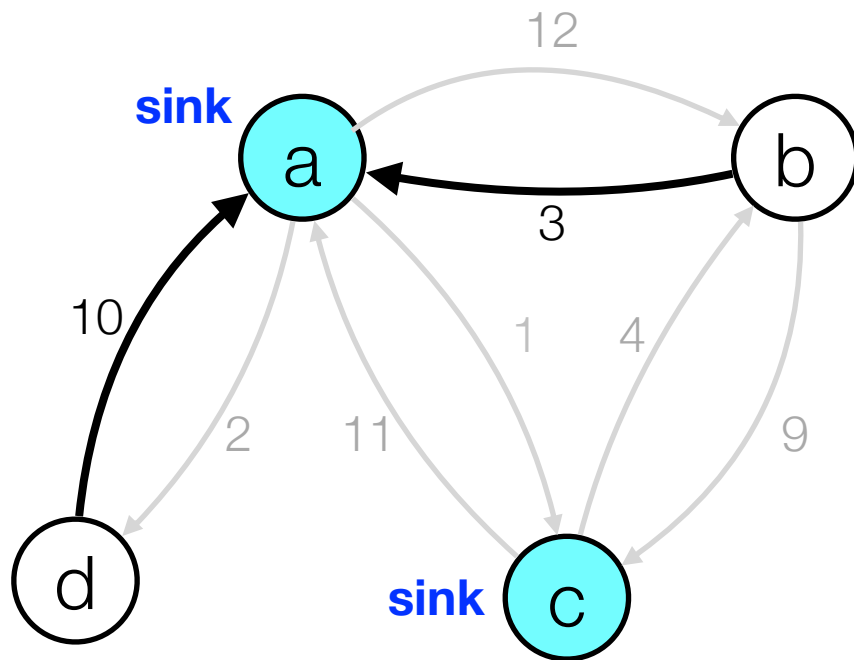
**Definition.** Let  $G(S,A,U)$  be a weighted directed graph.

A W-graph with  $k$  sinks is its subgraph satisfying:

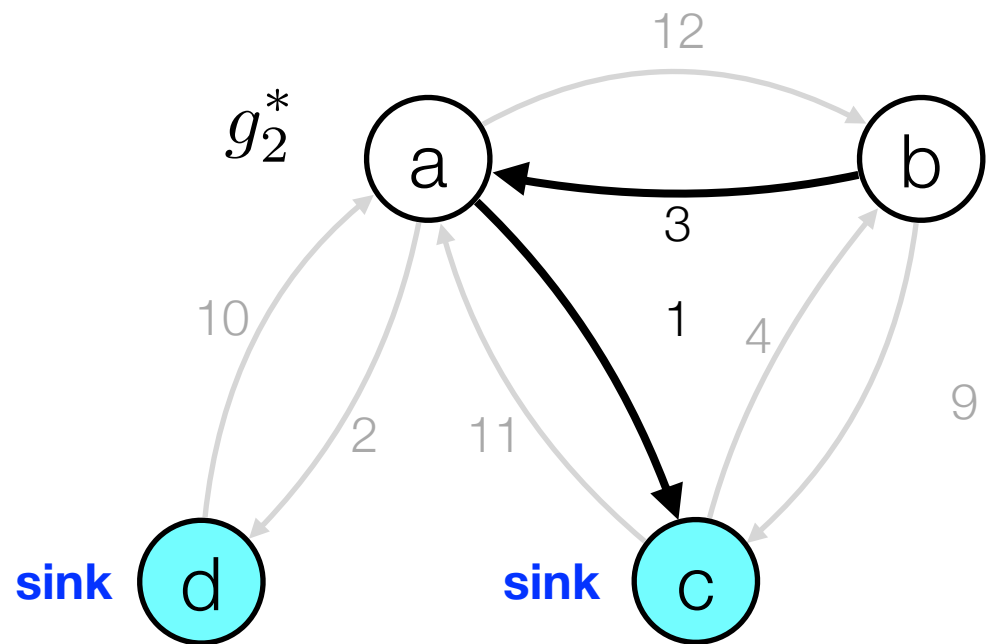
- (1) any sink has no outgoing arcs; any non-sink has exactly one outgoing arc;
- (2) the graph has no cycles.

## Example

A W-graph with two sinks



An optimal W-graph with two sinks

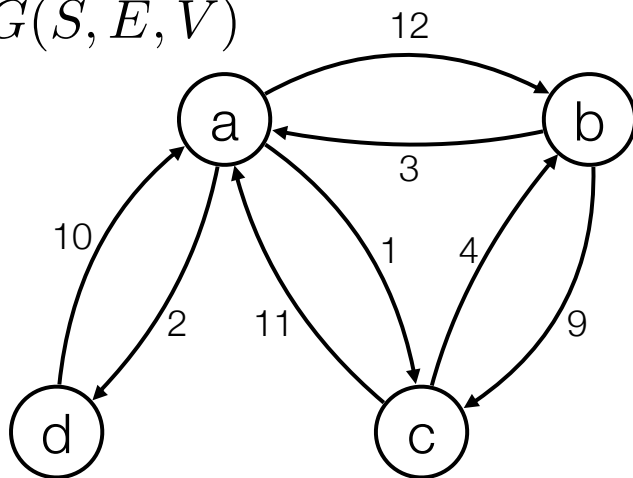


Optimal W-graphs with  $k$  sinks:

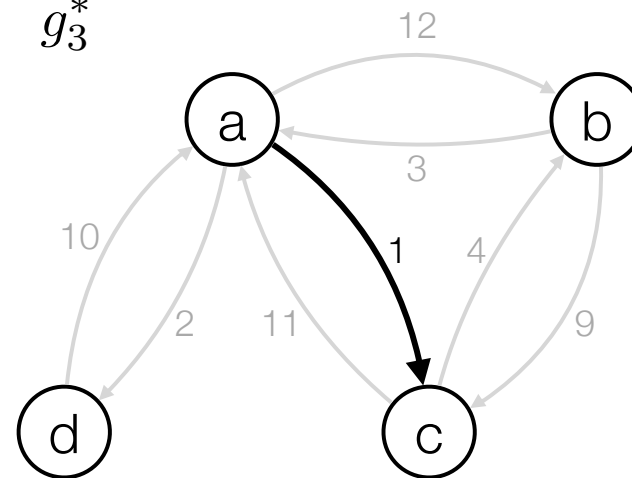
$\sum_{(i \rightarrow j) \in g} U_{ij}$  is minimized with respect to both  $k$  sinks and  $n-k$  arcs

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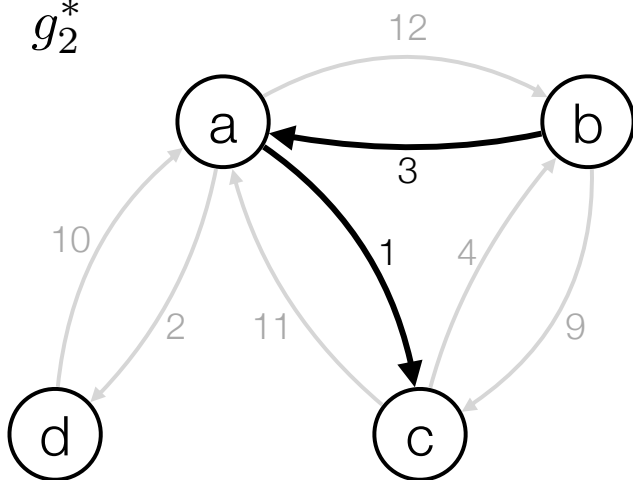
$G(S, E, V)$



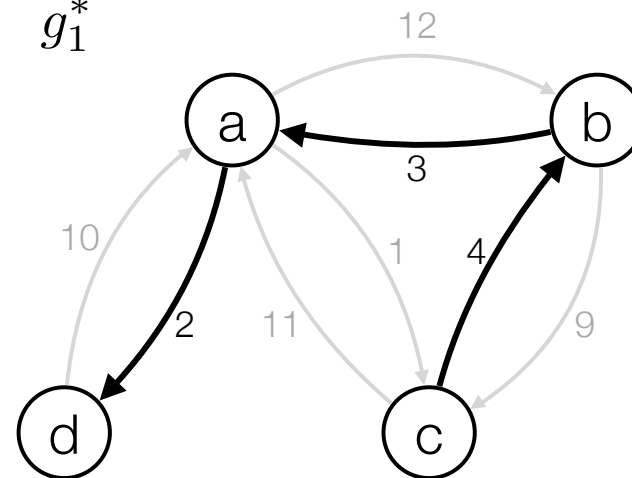
$g_3^*$



$g_2^*$

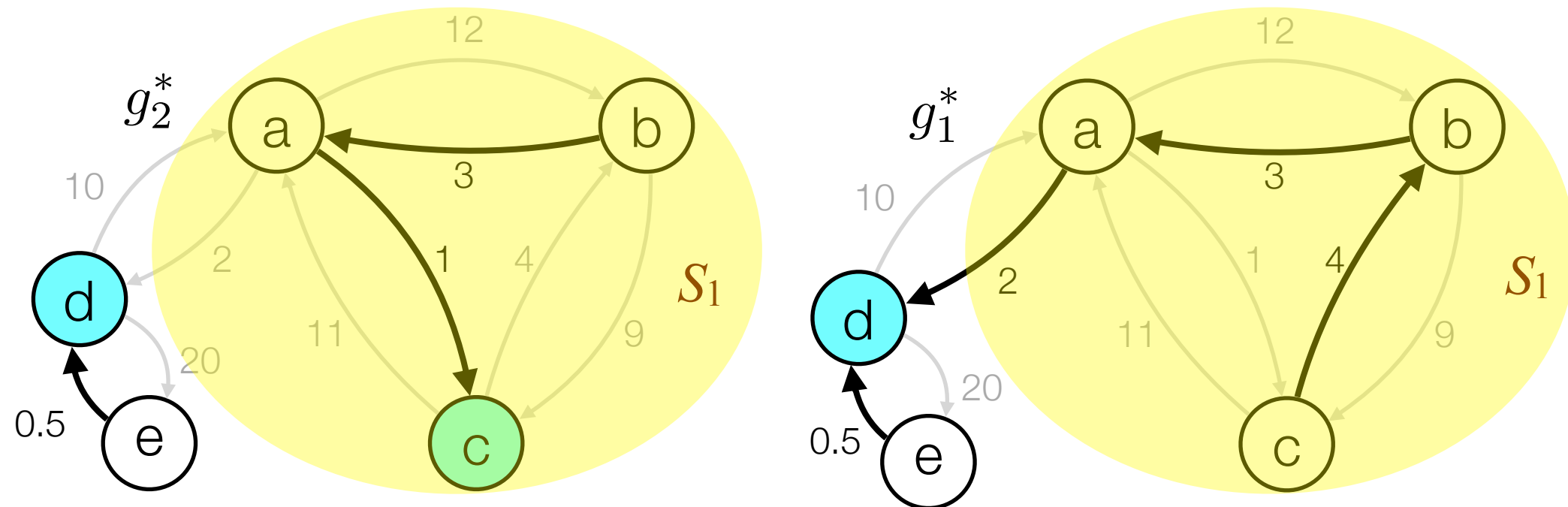


$g_1^*$



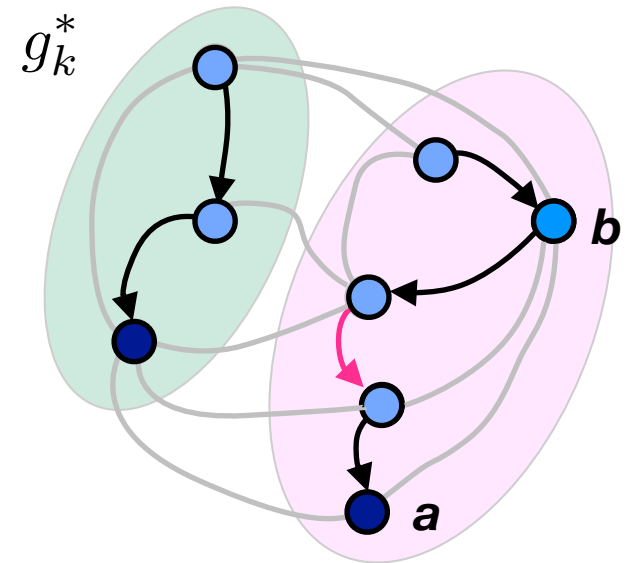
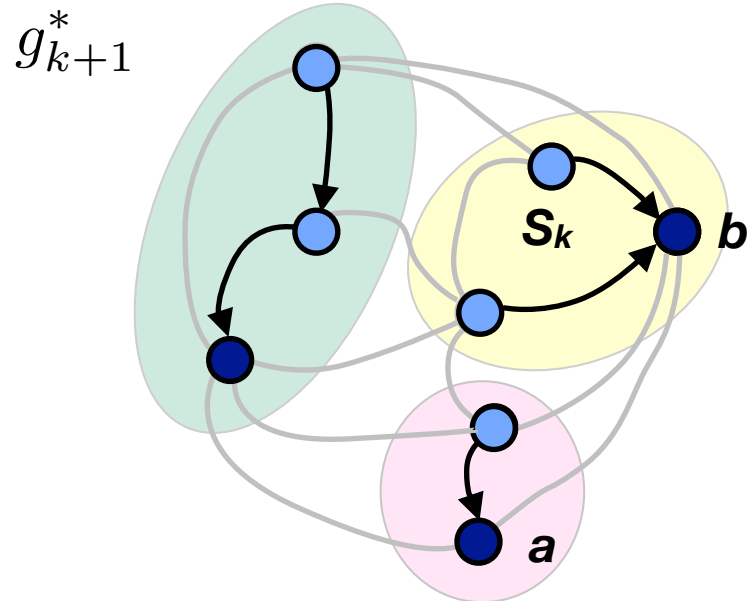
# Nested property (Gan, C., 2015)

- {The set of sinks of  $g_k^*$ }  $\subset$  {The set of sinks of  $g_{k+1}^*$ }
- There exists a connected component  $S_k$  of  $g_{k+1}^*$  whose set of vertices contains no sink of  $g_k^*$ .
- The sets of arcs connecting vertices  $S \setminus S_k$  in  $g_k^*$  and  $g_{k+1}^*$  coincide.
- In  $g_k^*$ , there is a single arc from  $S_k$  to  $S \setminus S_k$



# Asymptotics for eigenvectors

(under assumption that  $\parallel$  optimal W-graphs are unique)



Right eigenvectors:  $\phi_i^k = \begin{cases} 1, & i \in S_k \\ 0, & i \notin S_k \end{cases}$

**Time- reversible case:**  
Justification: Bovier, Eckort, Gayraud, Klein, early 2000's

Left eigenvectors:  $\psi_i^k = \begin{cases} 1, & i = b \\ -1, & i = a, \\ 0, & i \notin \{a, b\} \end{cases}$

**Time - irreversible case:**  
Justification:  
Cameron, Gan, 2015

Algorithms on graphs for finding asymptotic for  
eigenvectors and eigenvalues

# Algorithms for computing asymptotic spectra

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## **Algorithm 1**

(Cameron, 2014)

**Valid for time-reversible  
networks**

Precompute the  
minimum spanning tree

Order of computation:  
from smallest to largest

(Cameron, NHM, 9, 3, 383-416  
(2014), arXiv:1402.2869;  
Cameron. J. Chem. Phys., 141,

## *NEW!* **Algorithm 2**

(Tingyue Gan, Cameron 2015)

**Time-reversibility is  
not assumed**

A **single-sweep** algorithm  
(motivation: Chi-Liu/Edmond's algorithm  
for optimum branching for a directed  
graph with a selected root)

Order of computation:  
from largest to smallest

In preparation

# The single-sweep algorithm

(Cameron, Gan., 2015)

## Input:

list of arcs with weights  
and prefactors

$$1 \rightarrow i_1 : U_{1i_1}, k_{1i_1}$$

⋮

$$1 \rightarrow i_{n_1} : U_{1i_{n_1}}, k_{1i_{n_1}}$$

⋮

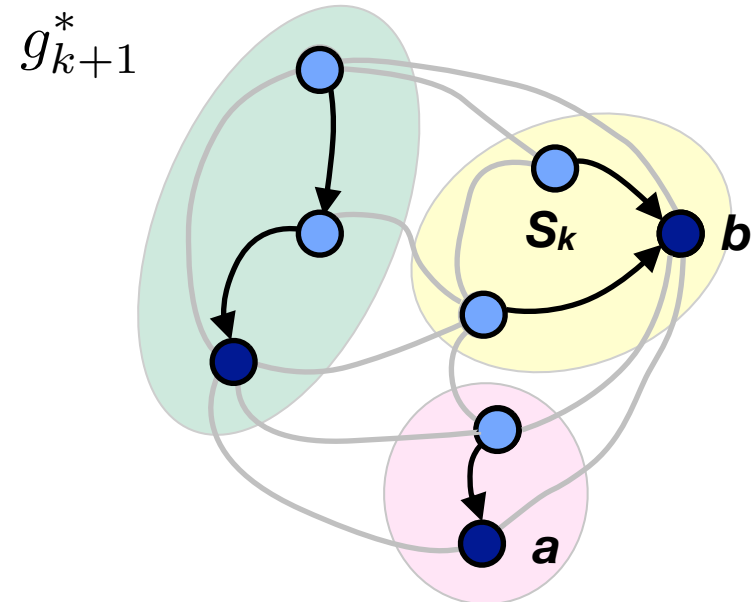
$$N \rightarrow p_1 : U_{1p_1}, k_{1p_1}$$

⋮

$$1 \rightarrow p_{n_N} : U_{1p_{n_N}}, k_{1p_{n_N}}$$

## Output:

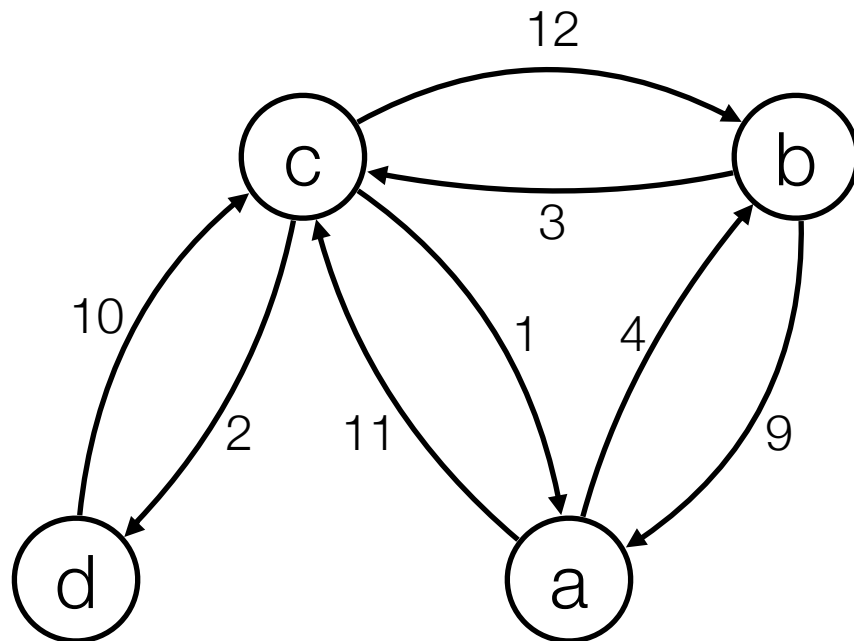
$$\begin{array}{cccccc} k = n - 1, & a_{n-1}, & \Delta_{n-1}, & A_{n-1}, & S_{n-1}, & b_{n-1} \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ k = 1, & a_1, & \Delta_1, & A_1, & S_1, & b_1 \end{array}$$



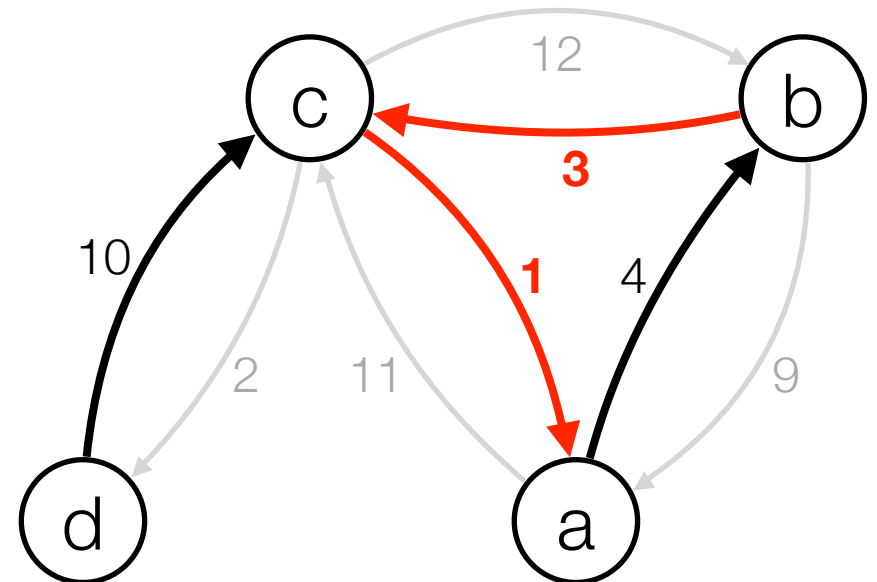
# The single-sweep algorithm

(T. Gan, C., 2015)

- Form binary trees out of sets of outgoing arcs for each vertex
- Delete the minimal outgoing arc  $\min\_arc(i)$  from the tree of each vertex  $i$  and add it to the main tree
- While the main tree contains more than 1 arc, keep adding arcs to build optimal W-graphs

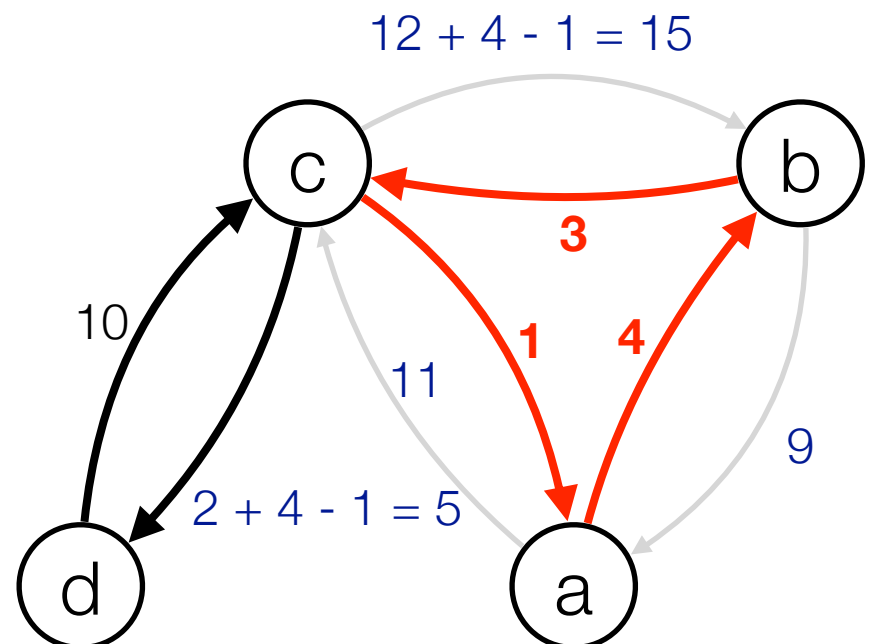
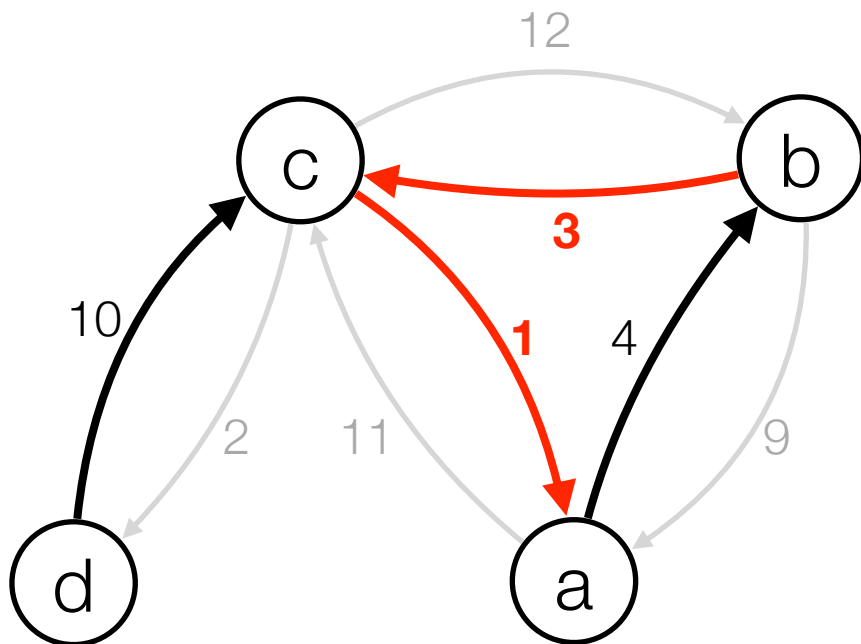


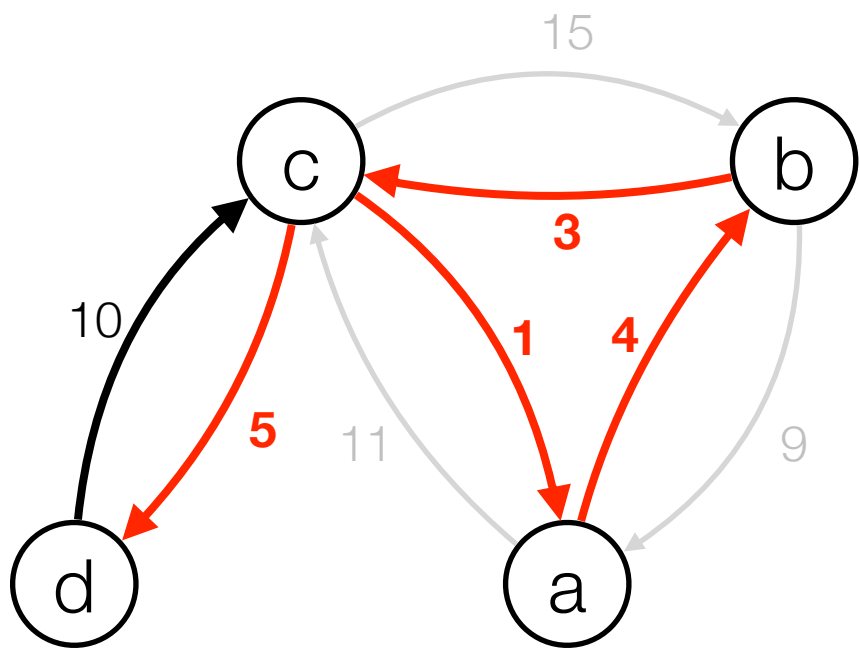
$$\Delta_3, = 1 \quad S_3 = \{c\}, \quad \text{sink}_3 = c$$
$$\Delta_2, = 3 \quad S_2 = \{b\}, \quad \text{sink}_2 = b$$





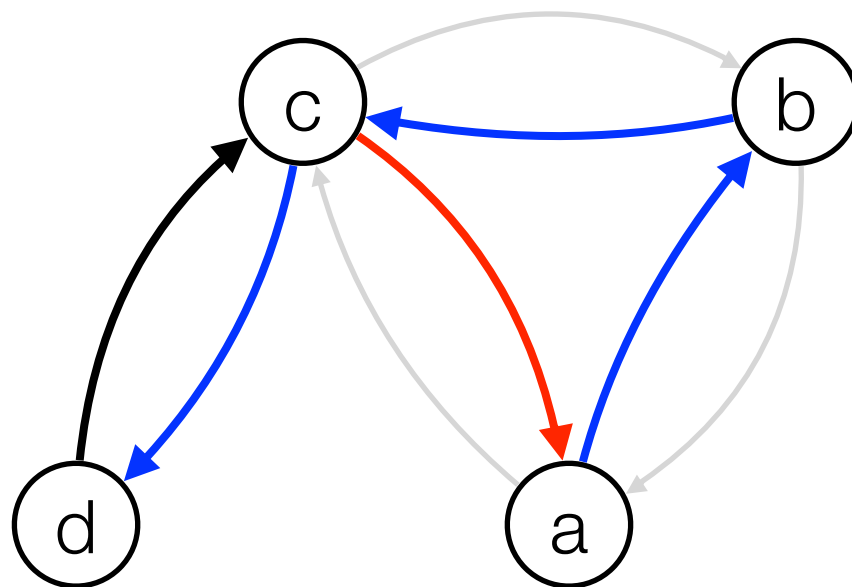
- If the arc  $a \rightarrow b$  on the top of the main tree creates a directed cycle, do:
  - Update  $U$ 's of in each tree reserve outgoing arcs of vertices lying
  - in the cycle according to the rule:  $U_{ij} = U_{ij} + U_{ab} - U_{\min\_arc(i)}$
  - Merge trees of remaining outgoing arcs from all vertices of the cycle
  - Keep deleting the minimal outgoing arcs  $p \rightarrow q$  from the merged tree and discard it until  $p$  and  $q$  are associated with the merged tree.
  - Delete the minimal outgoing arc  $p \rightarrow q$  from the merged tree and add it to the main tree.
  - For all vertices  $i$  associated with the merged tree,  $\min\_arc(i) = p \rightarrow q$





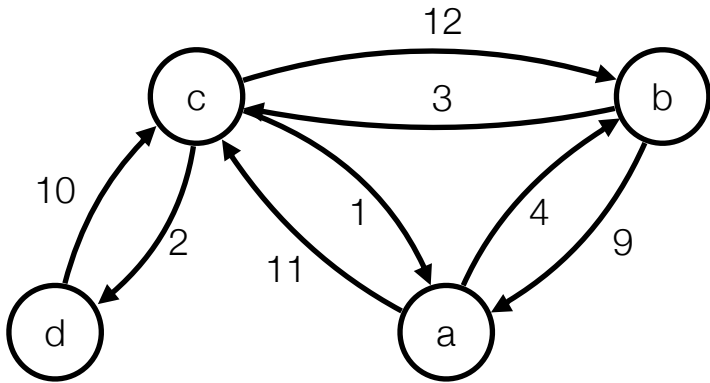
- Find optimal W-graphs:  
 recursively trace arcs from the set that ever appear on the top of the main tree backwards starting from the sink, so that each vertex has exactly one outgoing arc

$$\begin{aligned} \Delta_3, &= 1 & S_3 &= \{c\}, & \text{sink}_3 &= c \\ \Delta_2, &= 3 & S_2 &= \{b\}, & \text{sink}_2 &= b \\ \Delta_1, &= 5 & S_1 &= \{a, b, c\}, & \text{sink}_1 &= a \\ & & & & \text{sink}_0 &= d \end{aligned}$$

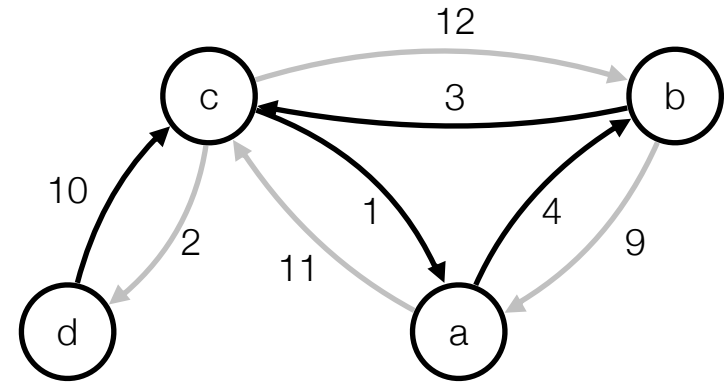


Blue arcs form the optimal W-graph  $g_1^*$

# Initialization

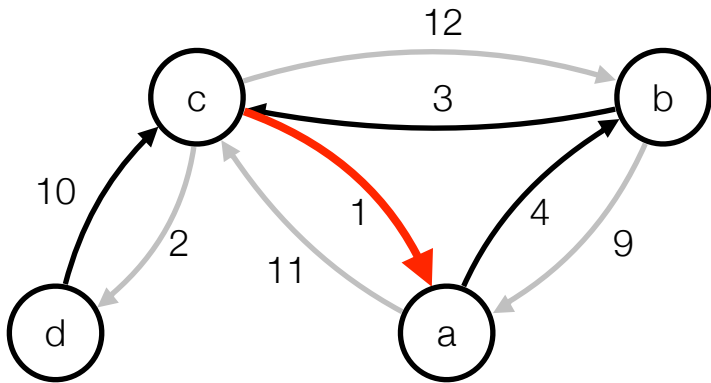


$$\begin{aligned}
 B_a &= \{[(a \rightarrow b), 4], [(a \rightarrow c), 11]\}; \\
 B_b &= \{[(b \rightarrow c), 3], [(b \rightarrow a), 9]\}; \\
 B_c &= \{[(c \rightarrow a), 1], [(c \rightarrow d), 2], [(c \rightarrow b), 12]\}; \\
 B_d &= \{[(d \rightarrow c), 10]\};
 \end{aligned}$$



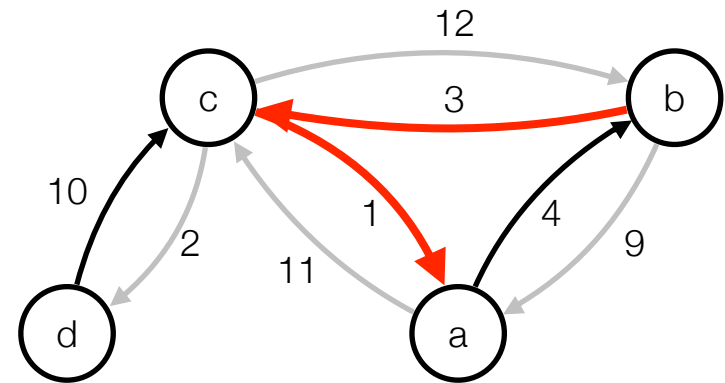
$$\begin{aligned}
 B_a &= \{[(a \rightarrow c), 11]\}; & B_b &= \{[(b \rightarrow a), 9]\}; \\
 B_c &= \{[(c \rightarrow d), 2], [(c \rightarrow b), 12]\}; & B_d &= \emptyset; \\
 M &= \{[(c \rightarrow a), 1], [(b \rightarrow c), 3], [(a \rightarrow b), 4], [(d \rightarrow c), 10]\};
 \end{aligned}$$

## While-cycle, step 1



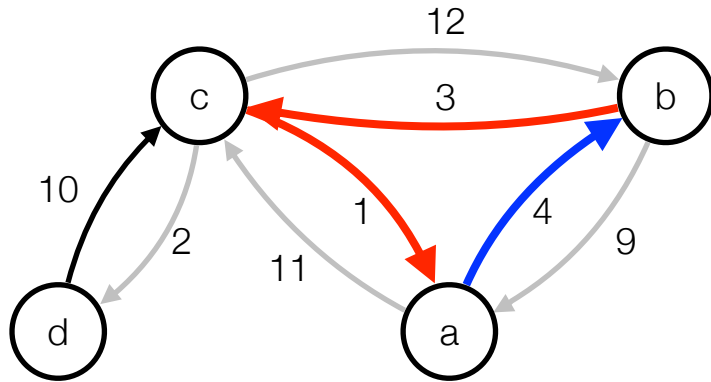
$$\begin{aligned}
 B_a &= \{[(a \rightarrow c), 11]\}; & B_b &= \{[(b \rightarrow a), 9]\}; \\
 B_c &= \{[(c \rightarrow d), 2], [(c \rightarrow b), 12]\}; & B_d &= \emptyset; \\
 M &= \{[(b \rightarrow c), 3], [(a \rightarrow b), 4], [(d \rightarrow c), 10]\}; \\
 s_3^* &= c; & \Delta_3 &= 1; & A_3 &= a_{ca}; \\
 g_3^* &= \{(c \rightarrow a)\};
 \end{aligned}$$

## While-cycle, step 2



$$\begin{aligned}
 B_a &= \{[(a \rightarrow c), 11]\}; & B_b &= \{[(b \rightarrow a), 9]\}; \\
 B_c &= \{[(c \rightarrow d), 2], [(c \rightarrow b), 12]\}; & B_d &= \emptyset; \\
 M &= \{[(a \rightarrow b), 4], [(d \rightarrow c), 10]\}; \\
 s_2^* &= b; & \Delta_2 &= 3; & A_2 &= a_{bc}; \\
 g_2^* &= \{(c \rightarrow a), (b \rightarrow c)\};
 \end{aligned}$$

## While-cycle, step 3



$$B_a = \{[(a \rightarrow c), 11]\}; \quad B_b = \{[(b \rightarrow a), 9]\};$$

$$B_c = \{[(c \rightarrow d), 2], [(c \rightarrow b), 12]\}; \quad B_d = \emptyset;$$

$$M = \{[(d \rightarrow c), 10]\};$$

Cycle  $\{(a \rightarrow b), (b \rightarrow c), (c \rightarrow a)\}$  is created

Update  $B_a$ ,  $B_b$ , and  $B_c$ :

$$B_a = \{[(a \rightarrow c), 11 + 4 - 4 = 11]\}; \quad B_b = \{[(b \rightarrow a), 9 + 4 - 3 = 10]\};$$

$$B_c = \{[(c \rightarrow d), 2 + 4 - 1 = 5], [(c \rightarrow b), 12 + 4 - 1 = 15]\};$$

Merge  $B_a$ ,  $B_b$ , and  $B_c$ :

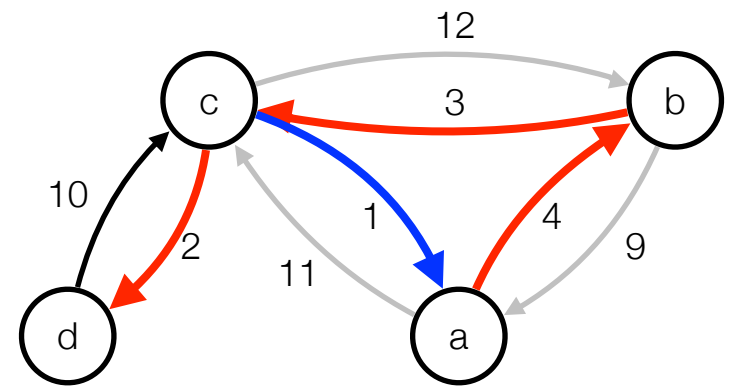
$$B := \{[(c \rightarrow d), 5], [(b \rightarrow a), 10], [(a \rightarrow c), 11], [(c \rightarrow b), 15]\};$$

$$B_a = B_b = B_c = B;$$

Remove the minimum arc from  $B$  and add it to  $M$ :

$$M = \{[(c \rightarrow d), 5], [(d \rightarrow c), 10]\};$$

## While-cycle, step 4



$$M = \{[(d \rightarrow c), 10]\};$$

$$s_1^* = a; \quad \Delta_1 = 5;$$

$$A_1 = \frac{a_{cd}a_{ab}}{a_{ca}};$$

$$g_1^* = \{(c \rightarrow d), (b \rightarrow c), (a \rightarrow b)\}$$

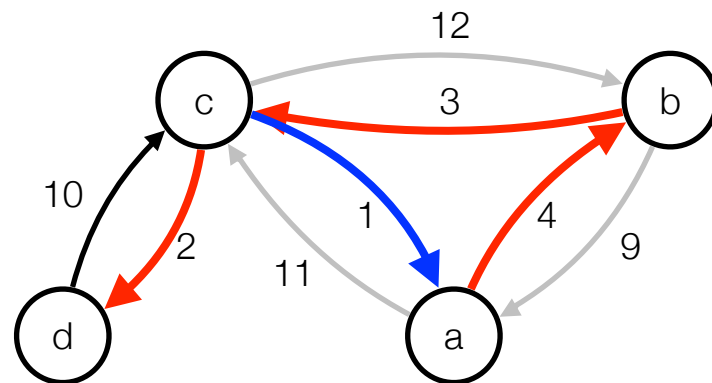
Results:

$$\lambda_3 \approx a_{ca} e^{-1/T}, \quad \phi_3 \approx \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}, \quad \psi_3 \approx [-1, 0, 1, 0]$$

$$\lambda_2 \approx a_{bc} e^{-3/T}, \quad \phi_2 \approx \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}, \quad \psi_2 \approx [-1, 1, 0, 0]$$

$$\lambda_1 \approx \frac{a_{cd} a_{ab}}{a_{ca}} e^{-5/T}, \quad \phi_1 \approx \begin{bmatrix} 1 \\ 1 \\ 1 \\ 0 \end{bmatrix}, \quad \psi_1 \approx [1, 0, 0, -1]$$

$$\lambda_0 = 0, \quad \phi_0 = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}, \quad \psi_0 \approx [0, 0, 0, 1]$$



# Computational cost

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$N$  vertices, index of each vertex  $\leq k$

**Best case scenario:**

Initialization:  $O(Nk \log k)$

Routine:  $O(N \log N)$

**Worst case scenario:**

Routine:  $O((Nk)^2 \log(Nk))$  due to merging trees of reserve arcs when a cycle is created

# Performance

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- **Lennard-Jones-38 network:**      **71887** vertices, **239706** arcs
  - CPU time: **30** seconds,
  - the number of cycles encountered: **50266**
  - the number of arcs having appeared on the top of the main tree: **122152**
- **Lennard-Jones-75 network:**      **169523** vertices, **441016** arcs
  - CPU time: **632** seconds (10.5 minutes)
  - the number of cycles encountered: **153164**
  - the number of arcs having appeared on the top of the main tree: **322686**

# Application to Lennard-Jones-75 network

Data: courtesy of David Wales

## Stats

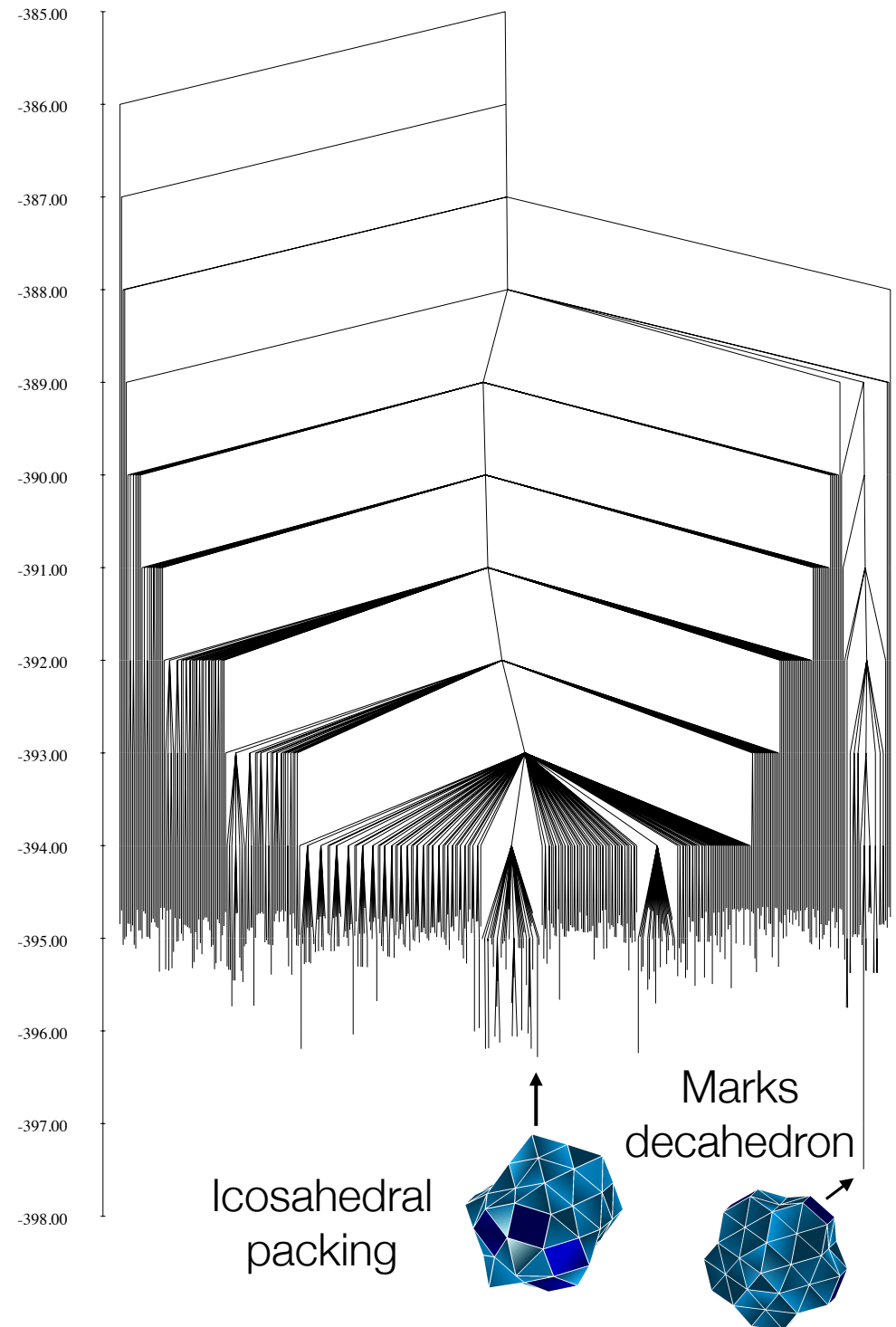
**593320** vertices, **452315** edges  
the maximal vertex index: **740**

## The maximal connected component:

**169523** vertices, **227198** edges  
the maximal vertex degree: **740**

the number of edges that are not loops and  
connecting different pairs of vertices:

**220508**

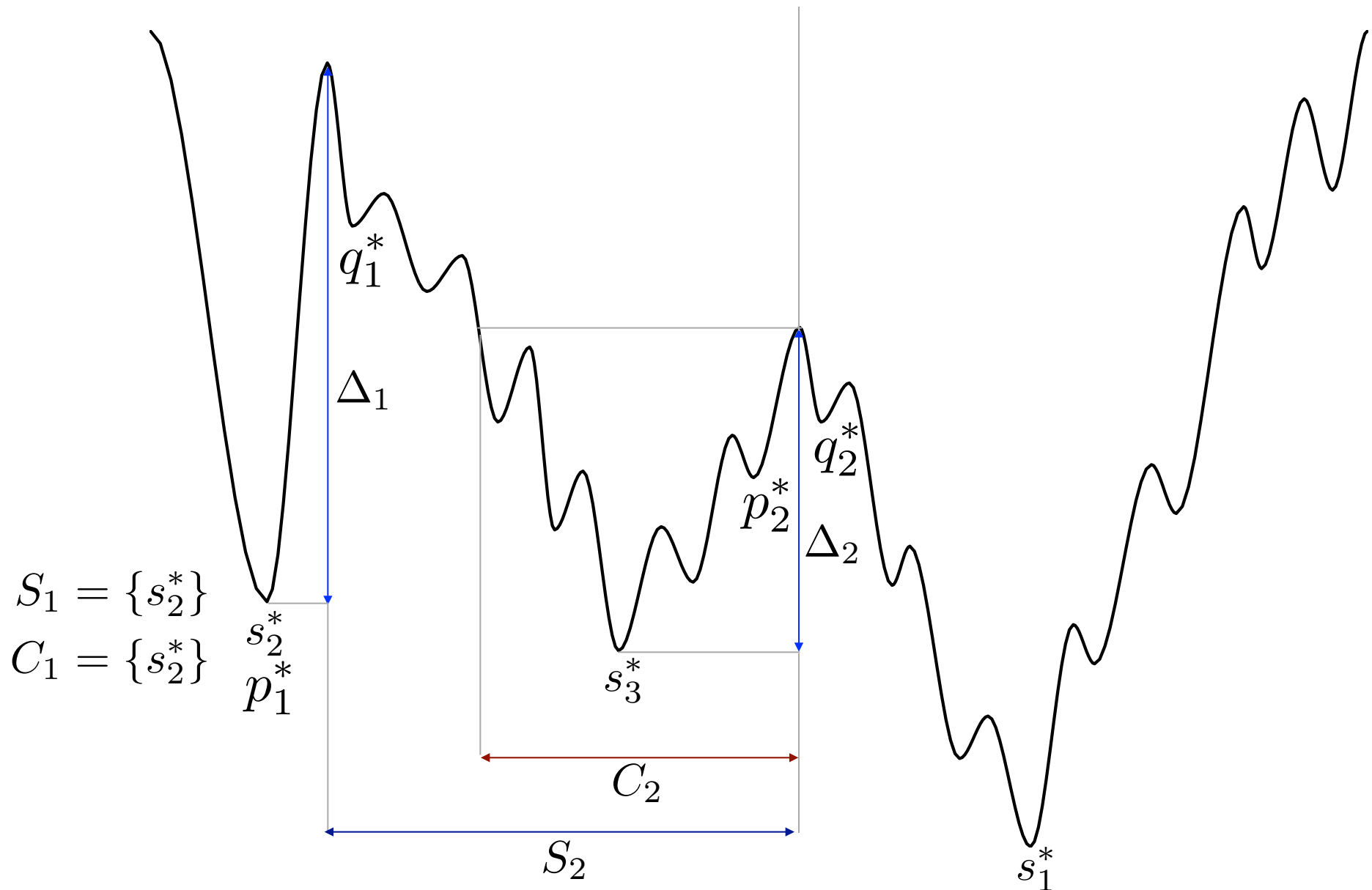




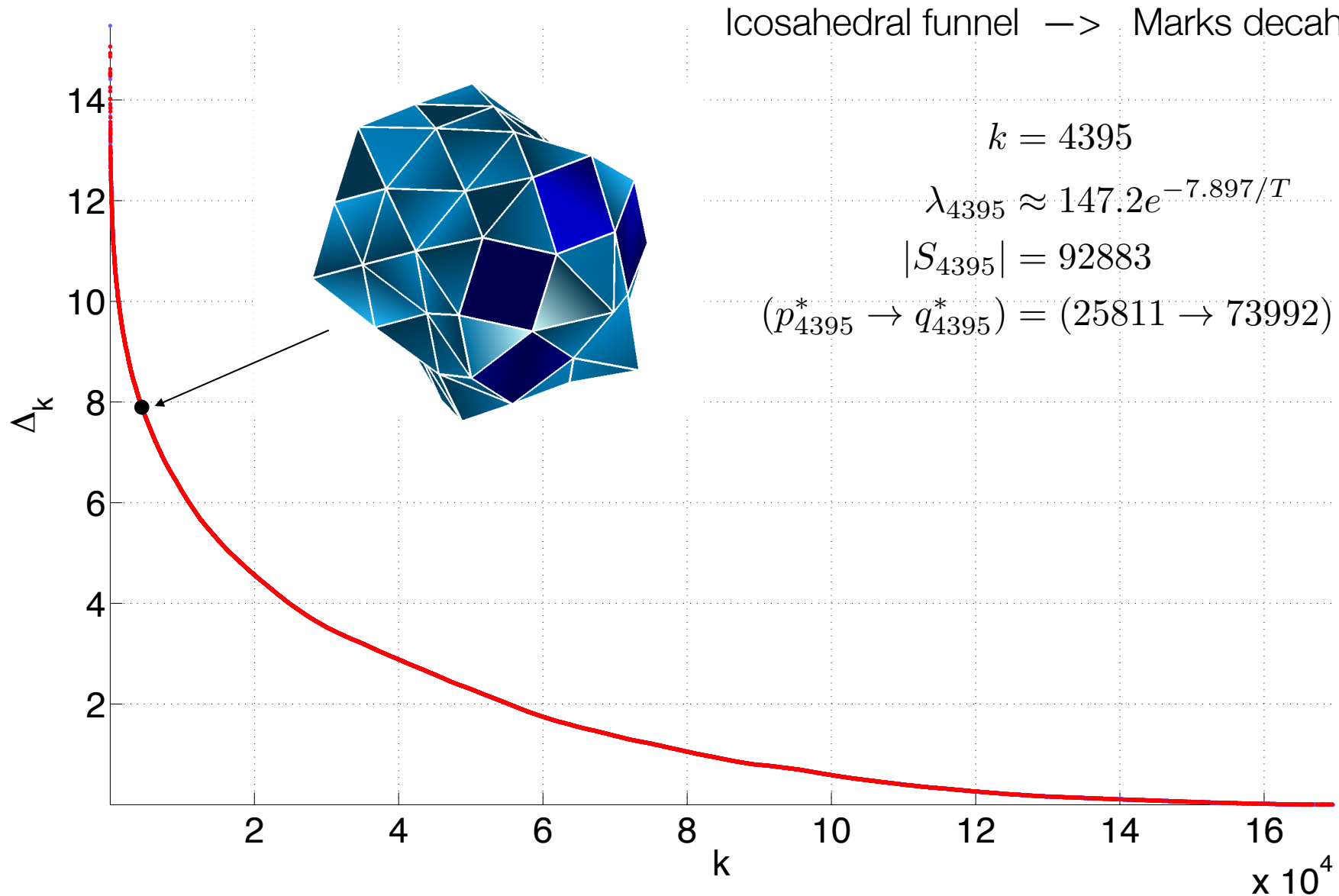
Zero-temperature asymptotic analysis to LJ<sub>75</sub>

# Quasi-invariant sets, Freidlin's cycles, etc.

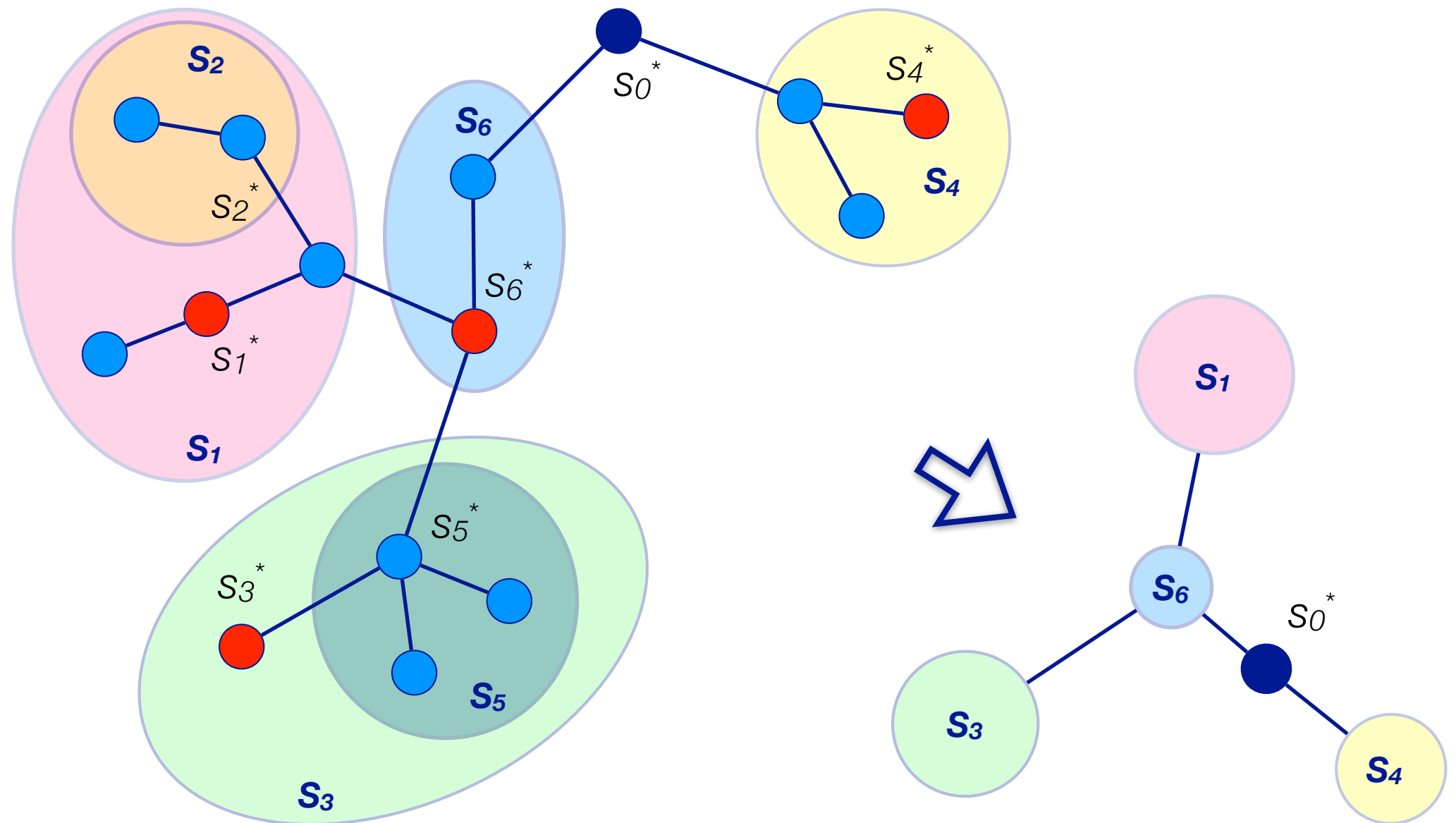
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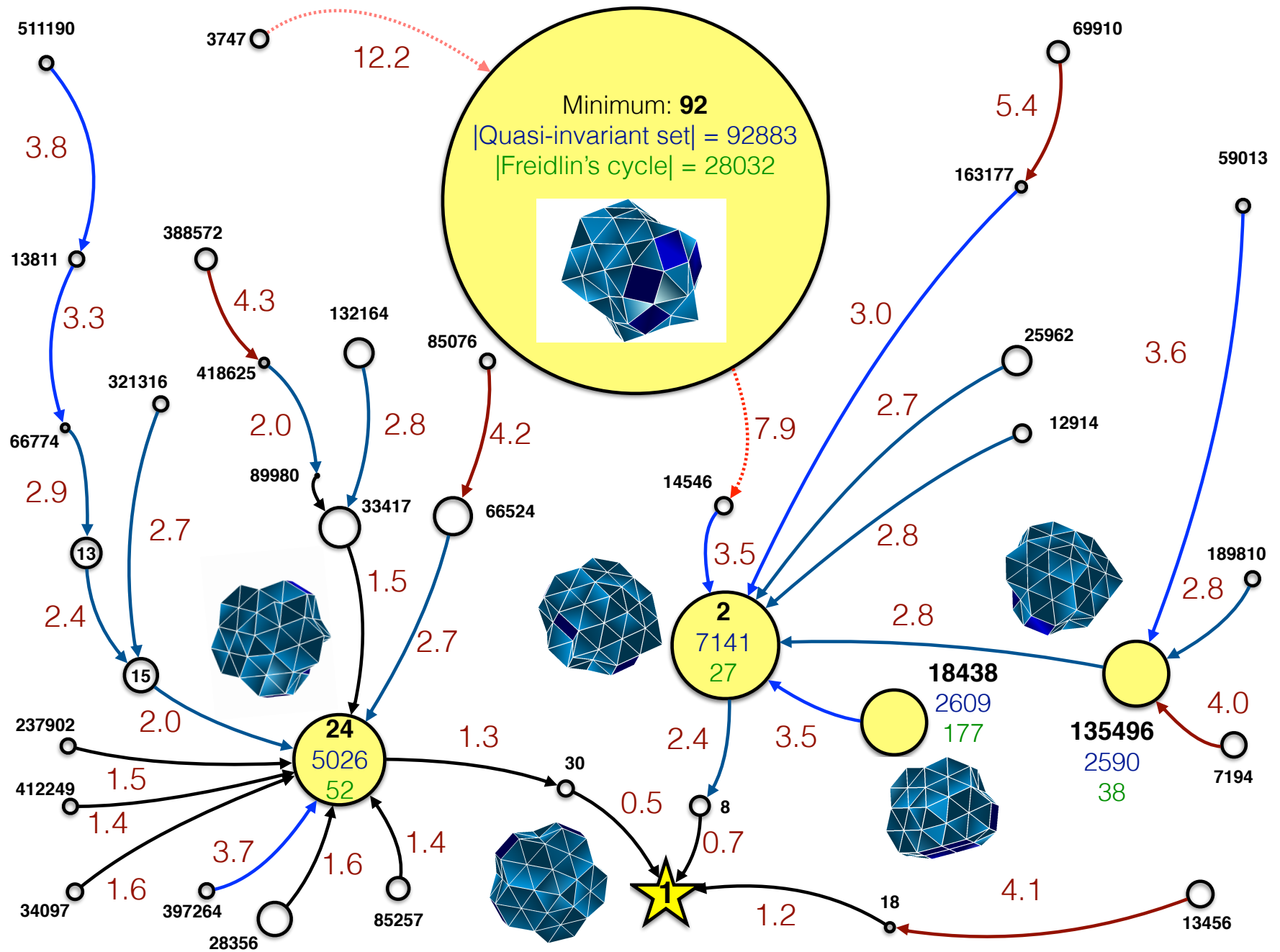
Asymptotics for eigenvalues:  $\lambda_k = \frac{O_{s_{k+1}^*} \bar{\nu}_{s_{k+1}^*}^{219}}{O_{p_k^* q_k^*} \bar{\nu}_{p_k^* q_k^*}^{218}} e^{-\Delta_k/T}$



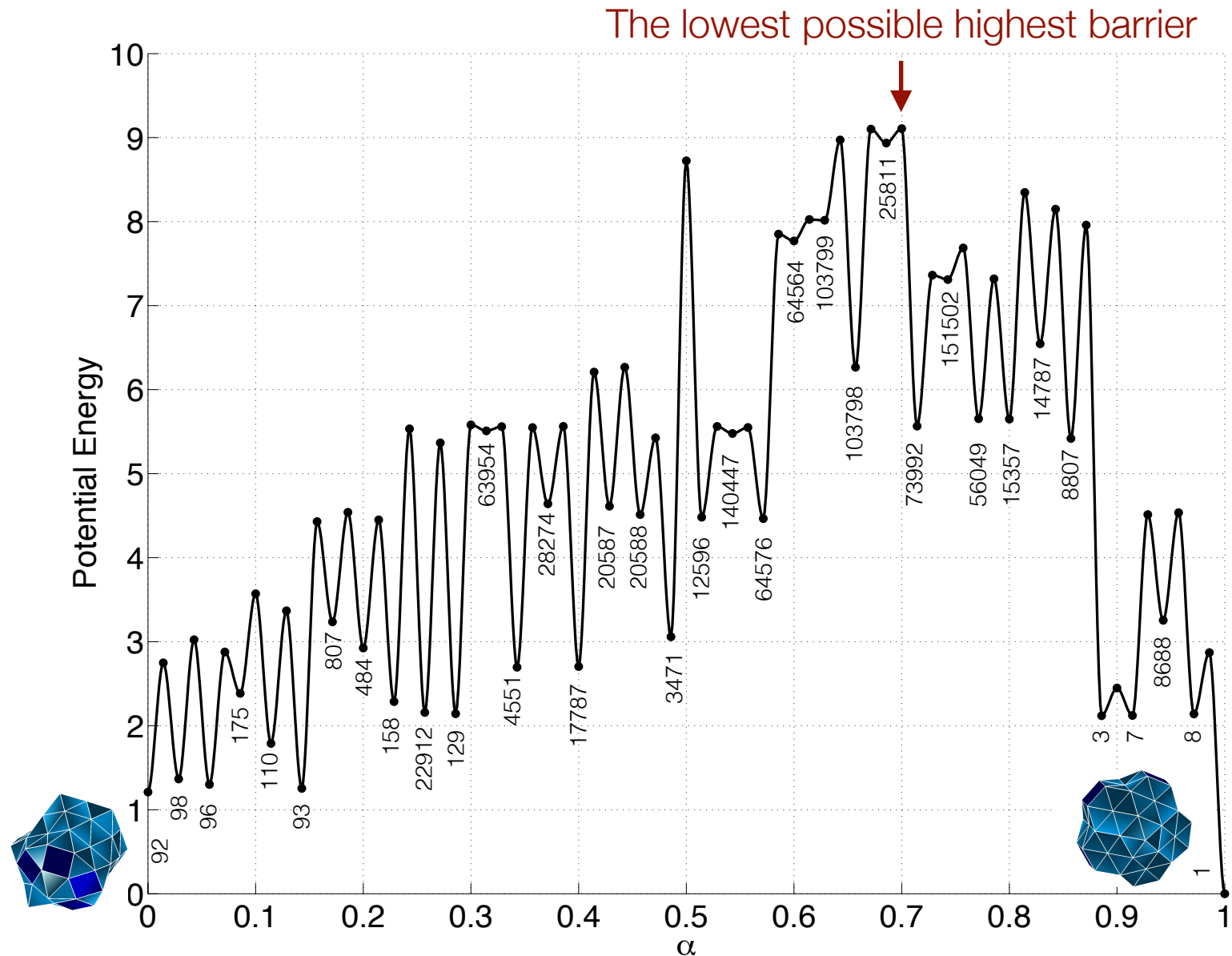
# Decomposition into maximal disjoint sets $S_k$



# Largest quasi-invariant sets ( $> 100$ local minima)



# The asymptotic zero-temperature (MinMax) path



# Crystalline Order Numbers (Steinhardt, Nelson, Ronchetti, 1983)

---

$$Q_l = \left[ \frac{4\pi}{2l+1} \sum_{m=-l}^l |\bar{Q}_{lm}|^2 \right]^{1/2}$$

Bond-orientational order numbers  
or crystalline order numbers

$$Q_{lm}(\mathbf{r}) \equiv Y_{lm}(\theta(\mathbf{r}), \phi(\mathbf{r}))$$

$Y_{lm}(\theta, \phi)$  Spherical harmonics

$\theta(\mathbf{r}), \phi(\mathbf{r})$  Polar angles of the bond  $\mathbf{r}$

$\bar{Q}_{lm}$  Averages over all bonds in the cluster where

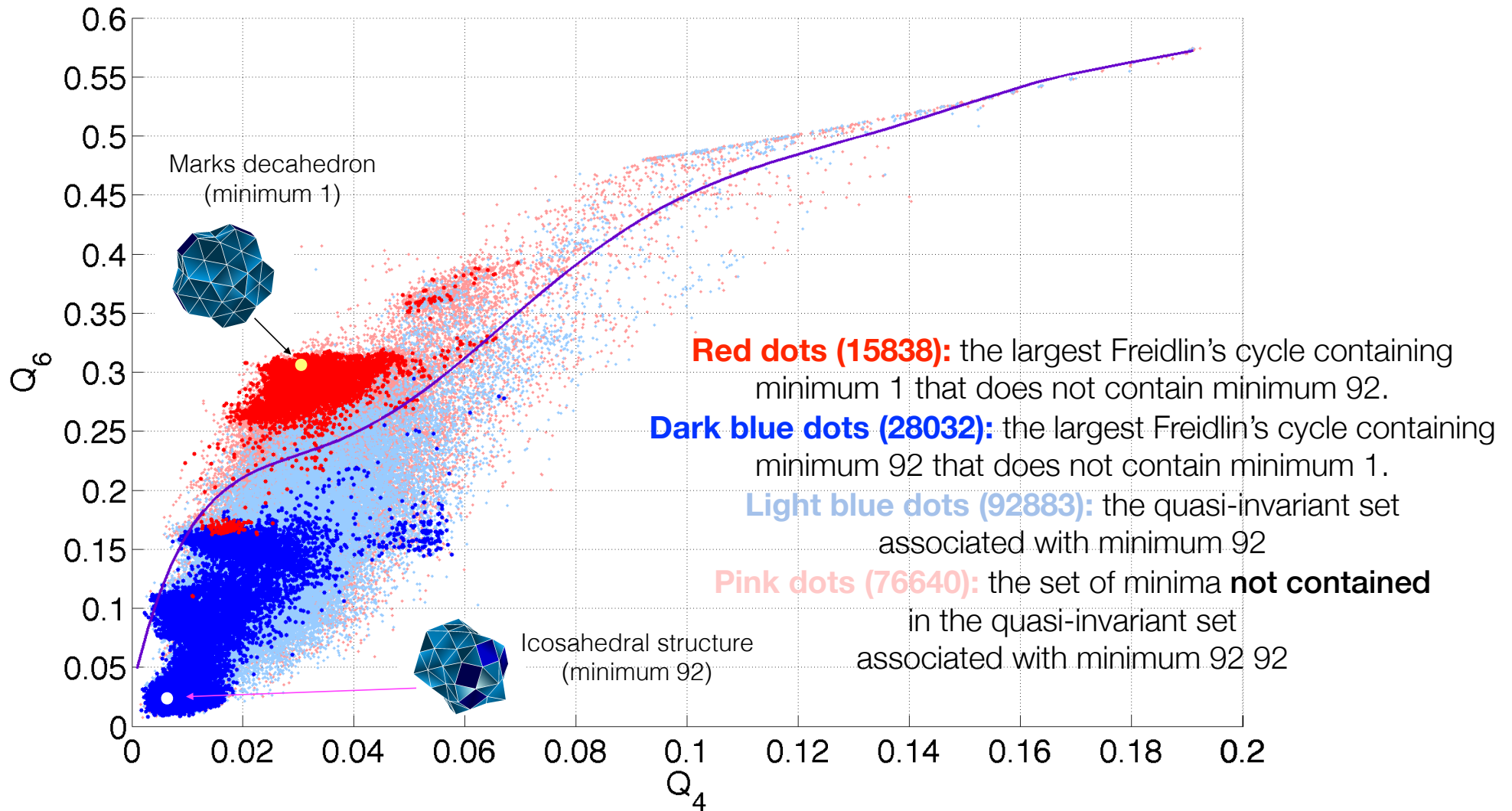
$$|\mathbf{r}| \leq 1.391r^*, \quad r^* = 2^{1/6}\epsilon$$

---

Following Picciani, Athenes, Kurchan, Tailleur, 2011, we use

$Q_4$  and  $Q_6$

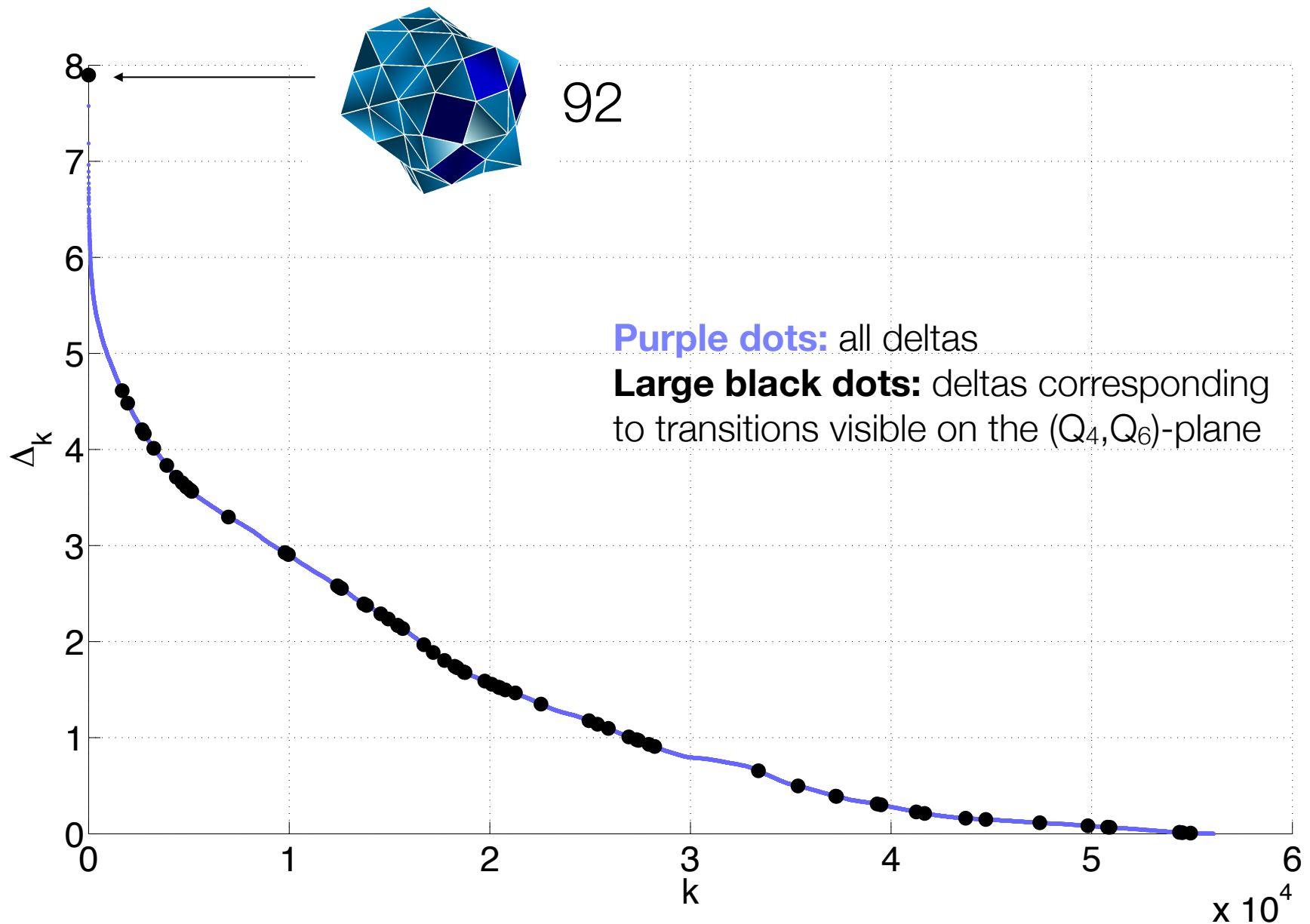
# Icosahedral and Marks decahedral basins



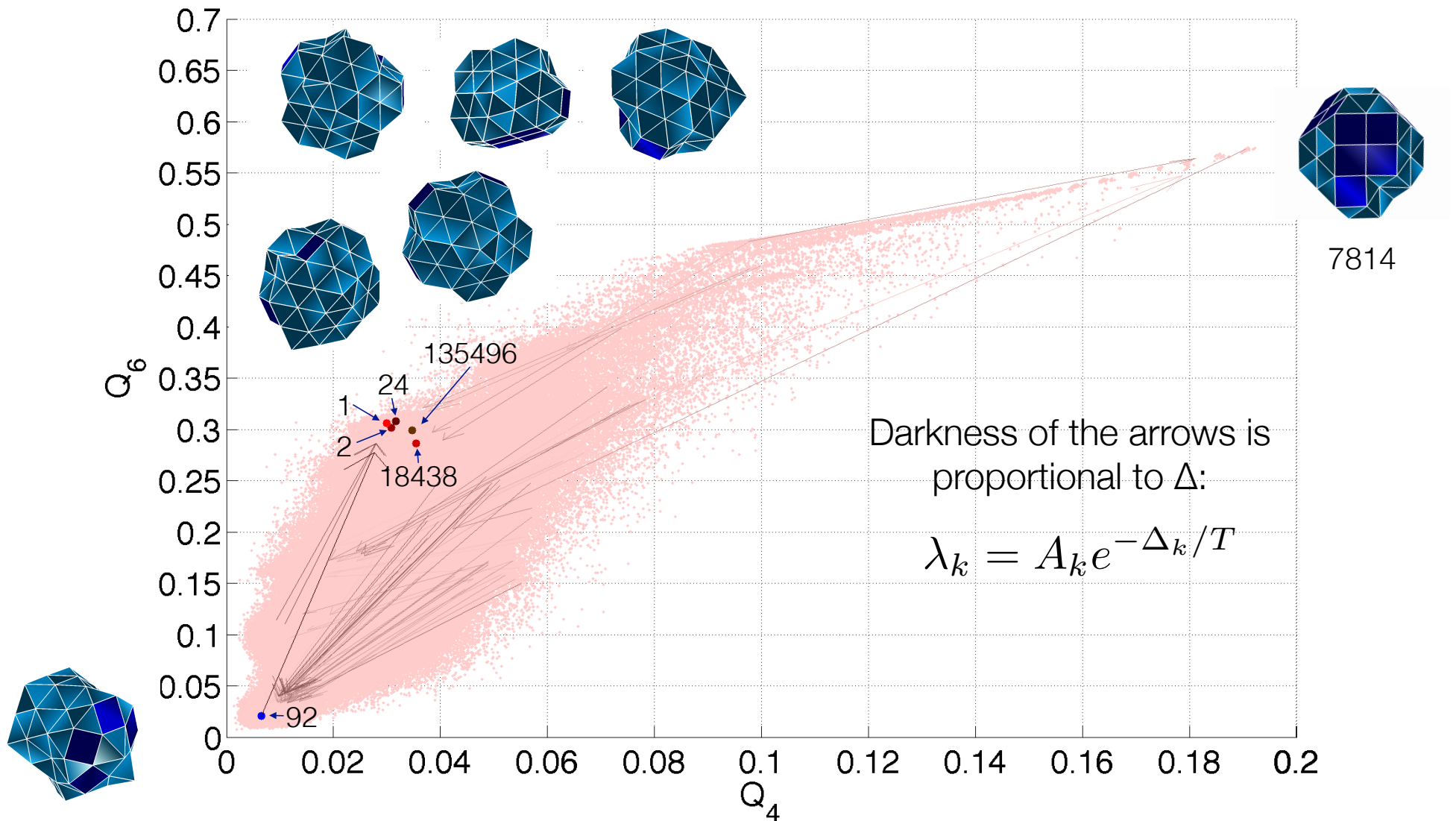


# Spectrum of the truncated LJ<sub>75</sub> network

( $V_{\max} - V_1 < 10.0$ ): **56074** vertices, **163666** arcs



# Relaxation processes visible on the $(Q_4, Q_6)$ -plane



Continuation to finite temperature of  $\lambda_{4395}$   
which is responsible for the relaxation process  
from the icosahedral funnel  
to Marks decahedron funnel

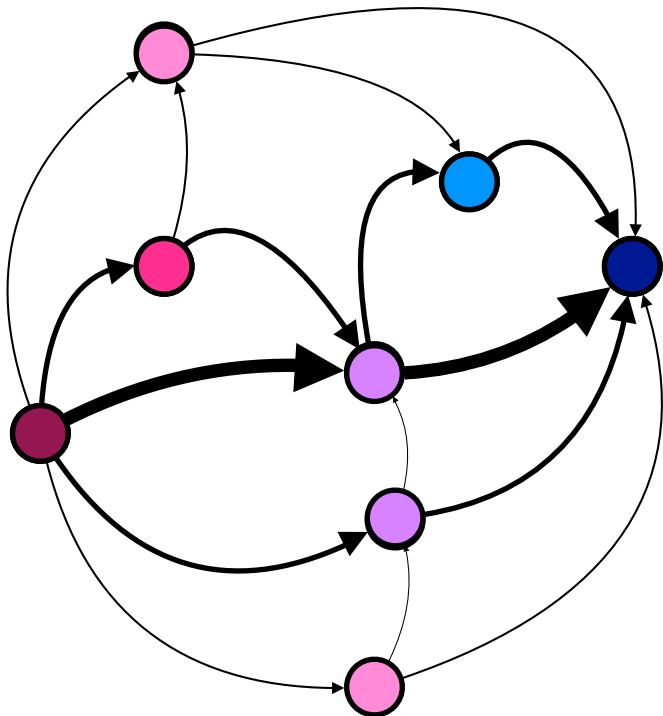
# Eigencurrent

$$F_{ij}^k := \pi_i L_{ij} e^{-\lambda_k t} [(\phi_k)_i - (\phi_k)_j]$$

The importance of currents was emphasized in works of J. Kurchan.

$\alpha_k F_{ij}^k$  = the net average number of transitions along the edge  $(i \rightarrow j)$  per unit time at time  $t$  in the relaxation process from the initial distribution  $\pi + \alpha_k P \phi_k$

E. Vanden-Eijnden proposed to consider eigencurrents.

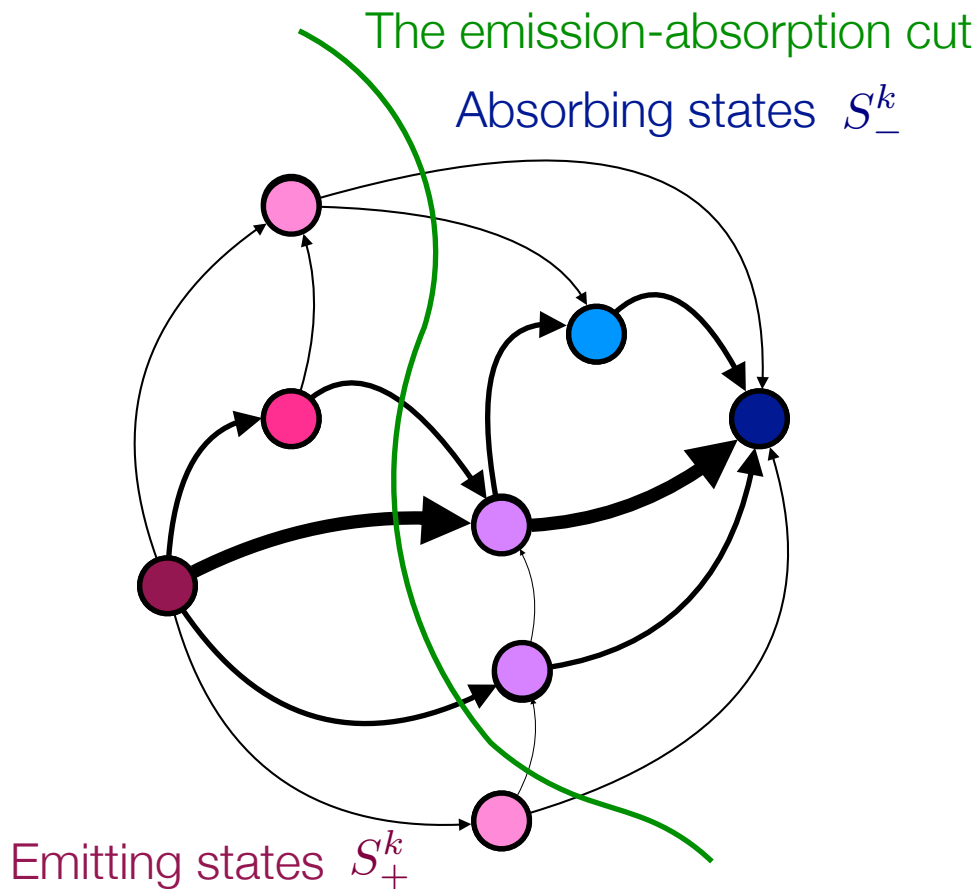


The Fokker-Planck equation in terms of eigencurrents

$$\frac{dp_i}{dt} = - \sum_{k=0}^{n-1} c_k \sum_{j \neq i} F_{ij}^k$$

$$\sum_{j \neq i} F_{ij}^k = e^{-\lambda_k t} \lambda_k \pi_k \phi_i^k$$

# The emission-absorption cut



Consider the total eigencurrent  $F^k$  through the vertex  $i$

$$\sum_{j \neq i} F_{ij}^k = e^{-\lambda_k t} \lambda_k \pi_k \phi_i^k$$

always  $> 0$     $> 0$  or  $< 0$

$$S = S_+^k \cup S_-^k$$

$$S_+^k := \{i \in S : (\phi_k)_i \geq 0\}$$

$$S_-^k := \{i \in S : (\phi_k)_i < 0\}$$

Among all possible cuts, the eigencurrent  $F^k$  is maximal through the emission-absorption cut

# Continuation of eigenpairs to finite temperatures

---

- **Difficulties:** (1) eigenvalues are close to 0 and may cross; (2) the matrix is large with widely varying entries
- **Useful fact:** the eigenvectors of the symmetrized generator matrix  $L_{sym} := P^{1/2}LP^{-1/2} \equiv P^{-1/2}QP^{-1/2}$  are orthonormal
- **Rayleigh Quotient iteration** with initial approximation
$$(\psi_k^0)_i = \begin{cases} \sqrt{\pi_i}, & i \in S_k \\ 0, & i \notin S_k \end{cases}$$
- **Precaution:** check whether the corresponding eigencurrent is largely emitted at the sink  $s_k^*$  and largely absorbed at the sink  $t_k^*$

# Rayleigh Quotient Iteration

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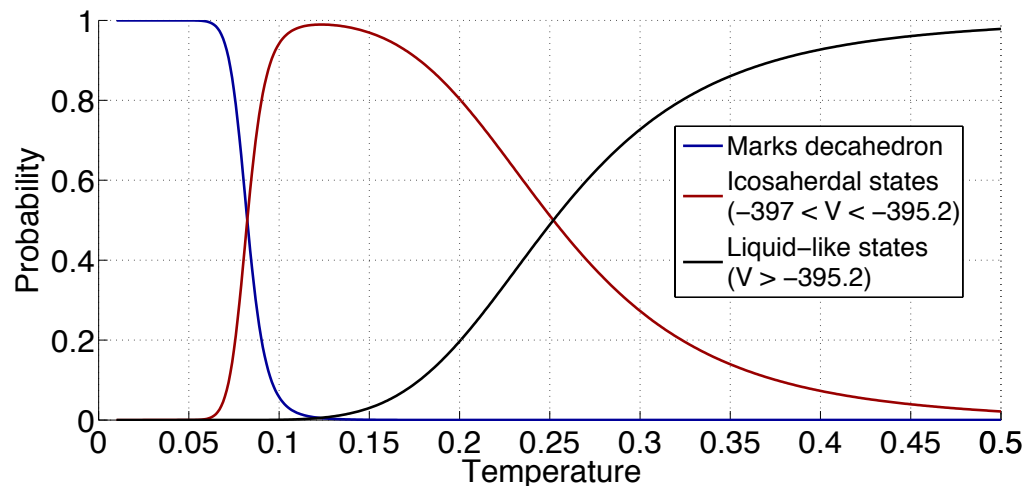
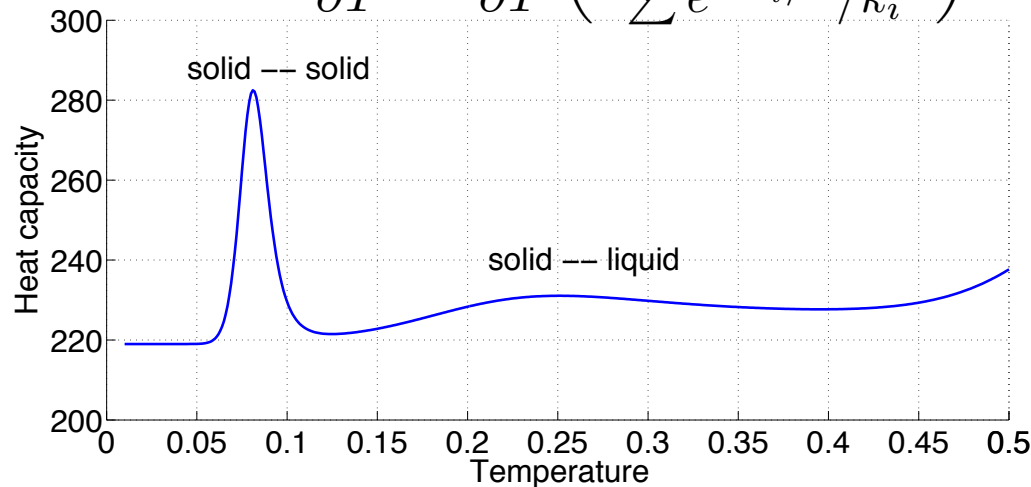
```
rayleigh = @(x) x'*Lsym*x/(x'*x); % the Rayleigh quotient

rtol = 1e-6;
itermax = 12;

%% Rayleigh quotient iteration
iter = 0;
while abs(res) > rtol*abs(lam) & iter < itermax
    A = Lsym - lam*speye(n);
    w = (A)\v;
    v = w/norm(w);
    res = norm(A*v);
    lam = rayleigh(v);
    iter = iter + 1;
    fprintf('iter = %d:  lam = %d\t res = %d\n',iter,lam,res);
end
```

# Difficulties with Lennard-Jones-75

$$c_v = \frac{\partial \langle E \rangle}{\partial T} = \frac{\partial}{\partial T} \left( \frac{\sum E_i e^{-E_i/T} / k_i}{\sum e^{-E_i/T} / k_i} \right)$$



Marks decahedron - icosahedral states  
solid - solid transition:  $T = 0.08$

Icosahedral - liquid-like states  
transition:  $T = 0.25$

The range of temperatures to which  
we would like to continue  $\lambda_{4395}$  :

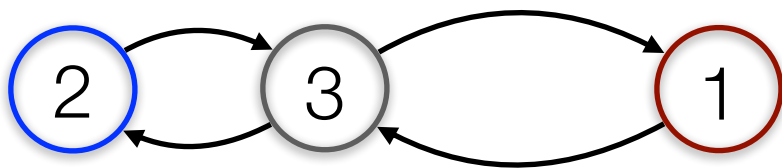
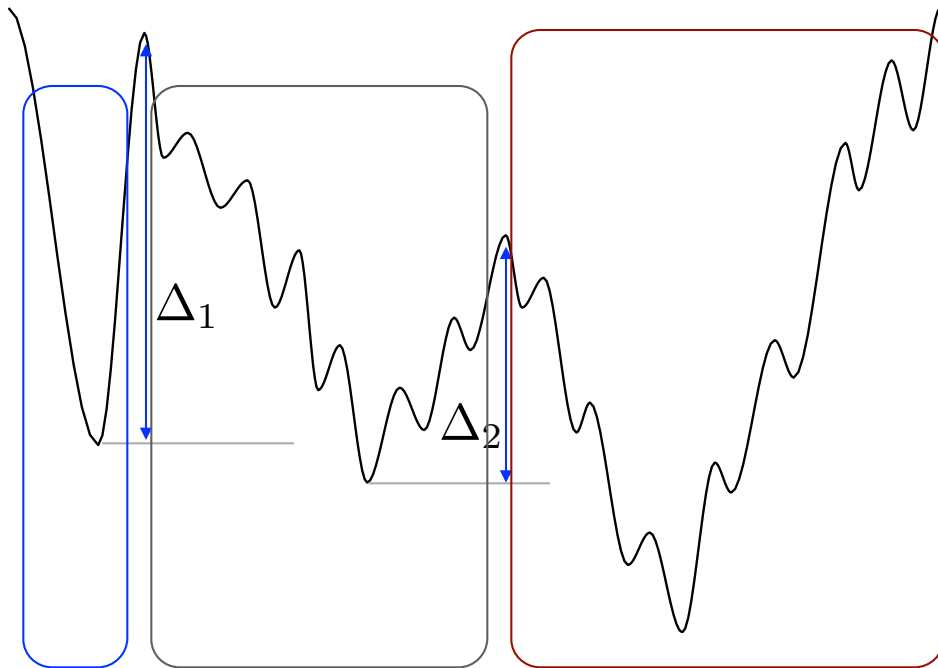
$$0.05 \leq T \leq 0.25$$

For  $T < 0.17$ , the matrix is badly  
scaled, and the results are  
inaccurate or NaN

For  $T \geq 0.17$ , convergence  
to a wrong eigenpair takes place



# Remedy 1: lumping



The lumped network

Pick  $\Delta_{\min}$ . Here  $\Delta_{\min} = \Delta_2$

Lump the quasi-invariant sets with  $\Delta_k < \Delta_{\min}$

Re-calculate pairwise rates  $S_k$

$$\tilde{L}_{kl} = \sum_{i \in S_k, j \in S_l} L_{ij} \frac{\pi_i}{\sum_{i' \in S_k} \pi_{i'}}$$

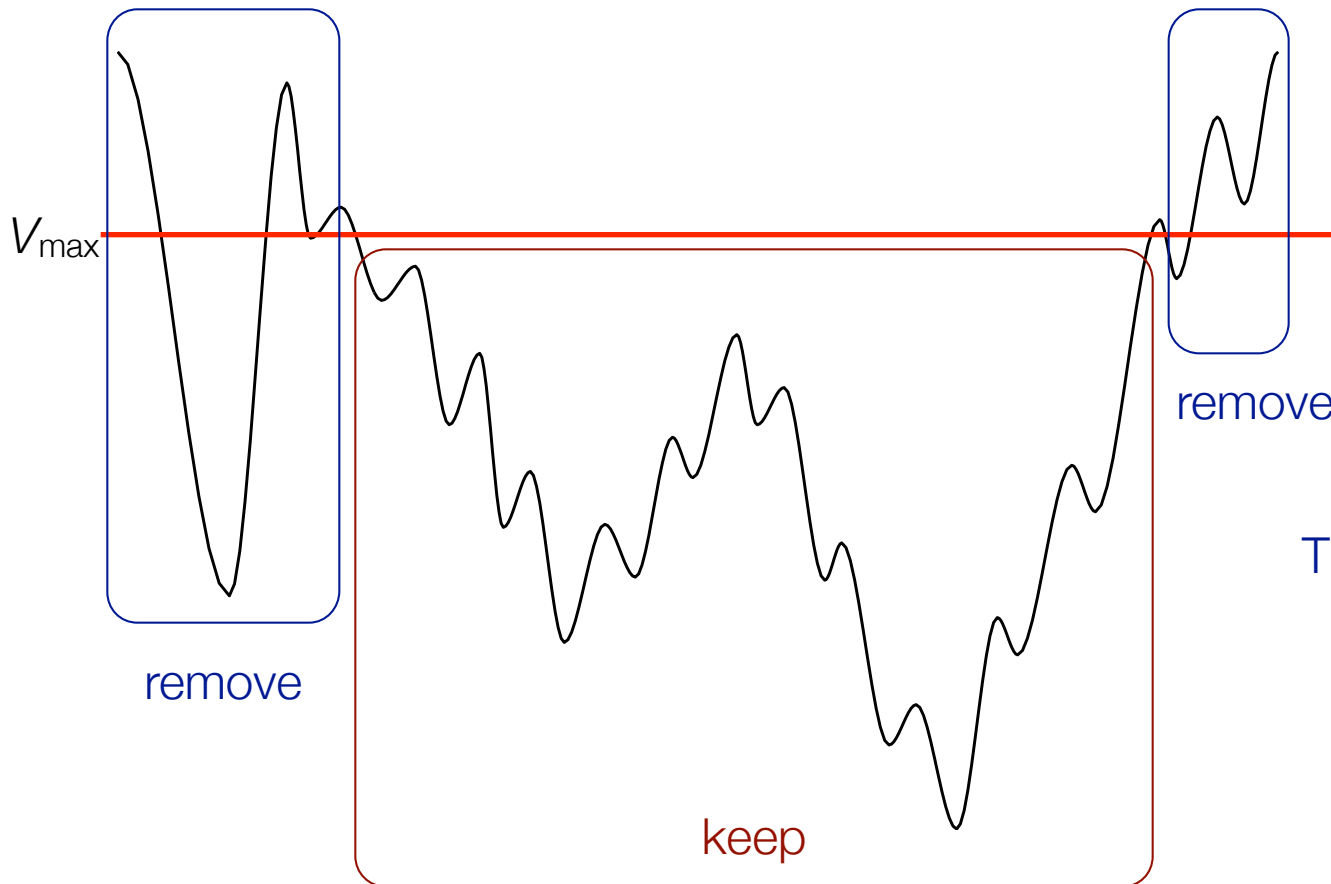
The resulting generator matrix  $\tilde{L}$  is smaller, the largest entries of  $L$  are gone

$$\tilde{L}_{N \times N} = \begin{bmatrix} A_{N \times n} & L_{n \times n} & B_{n \times N} \end{bmatrix}$$

$$A_{ki} = \begin{cases} \frac{\pi_i}{\sum_{i' \in S_k} \pi_{i'}}, & i \in S_k \\ 0, & \text{otherwise} \end{cases}$$

$$B_{jl} = \begin{cases} 1, & j \in S_l \\ 0, & \text{otherwise} \end{cases}$$

# Remedy 2: truncation

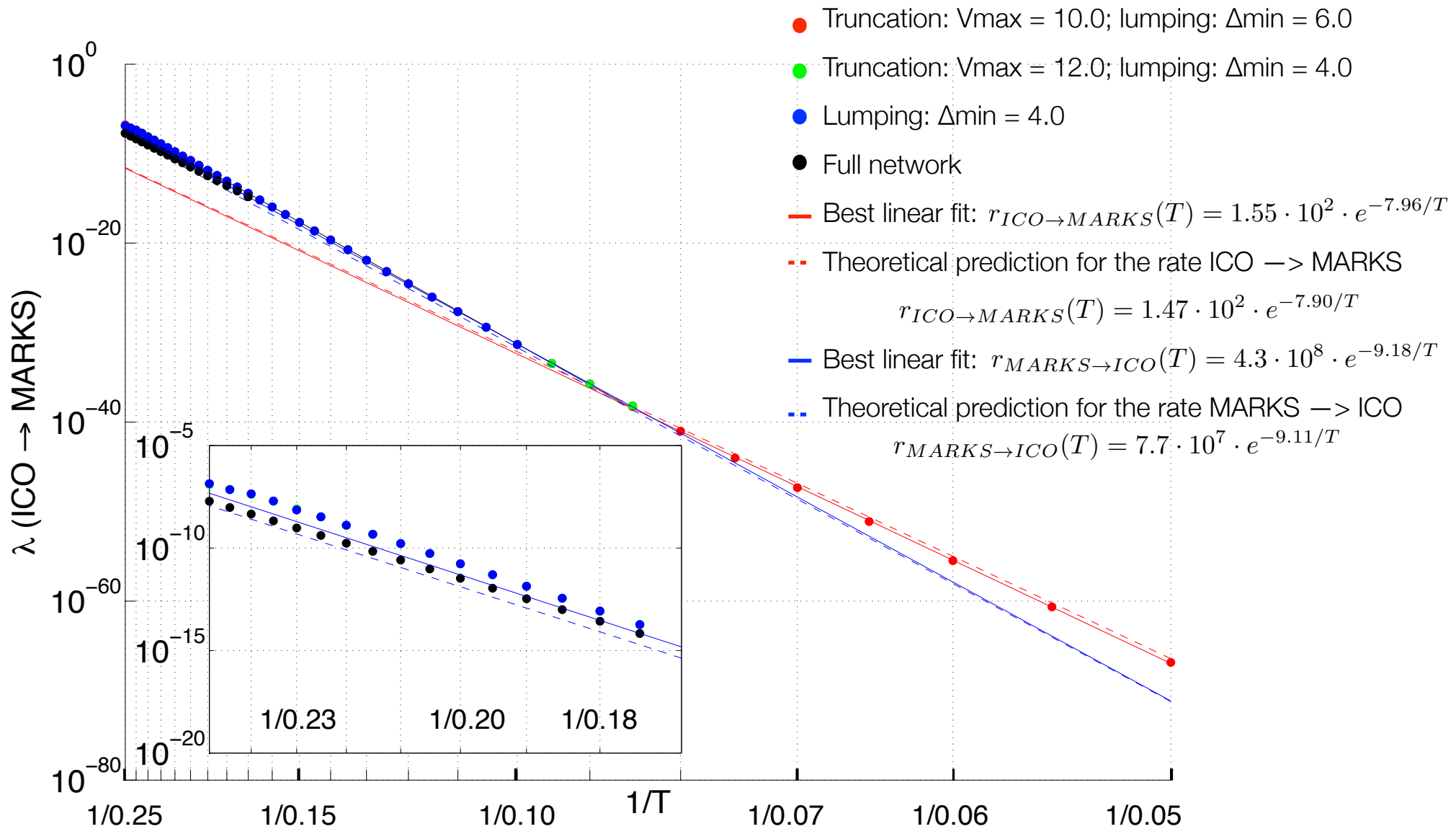


Pick  $V_{\max}$ , remove all states separated from the global minimum by a barrier exceeding  $V_{\max}$

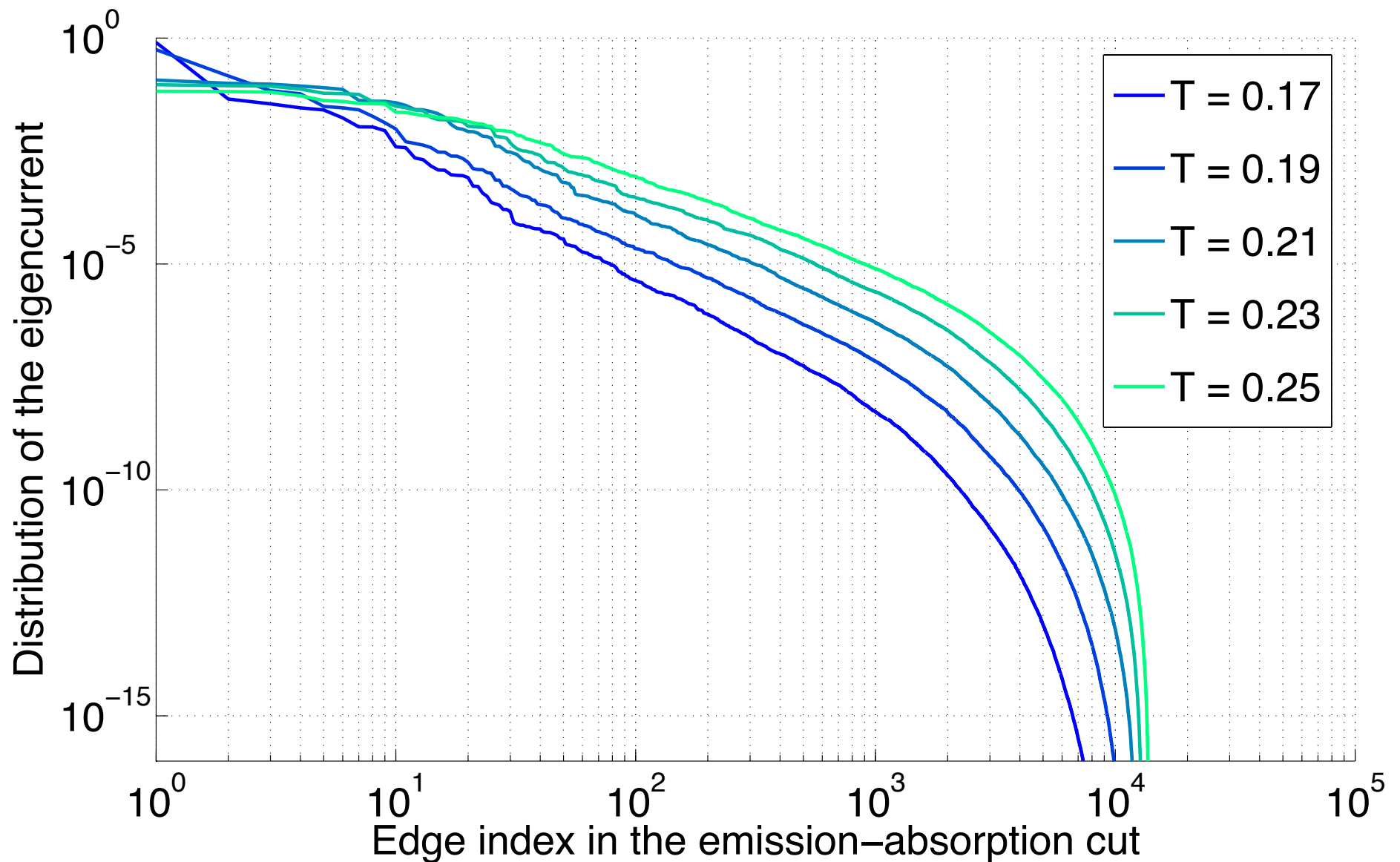
The resulting network is smaller, the components that used to be nearly transient or make it nearly reducible are removed

One can combine truncation and lumping:  
first truncate, then lump.

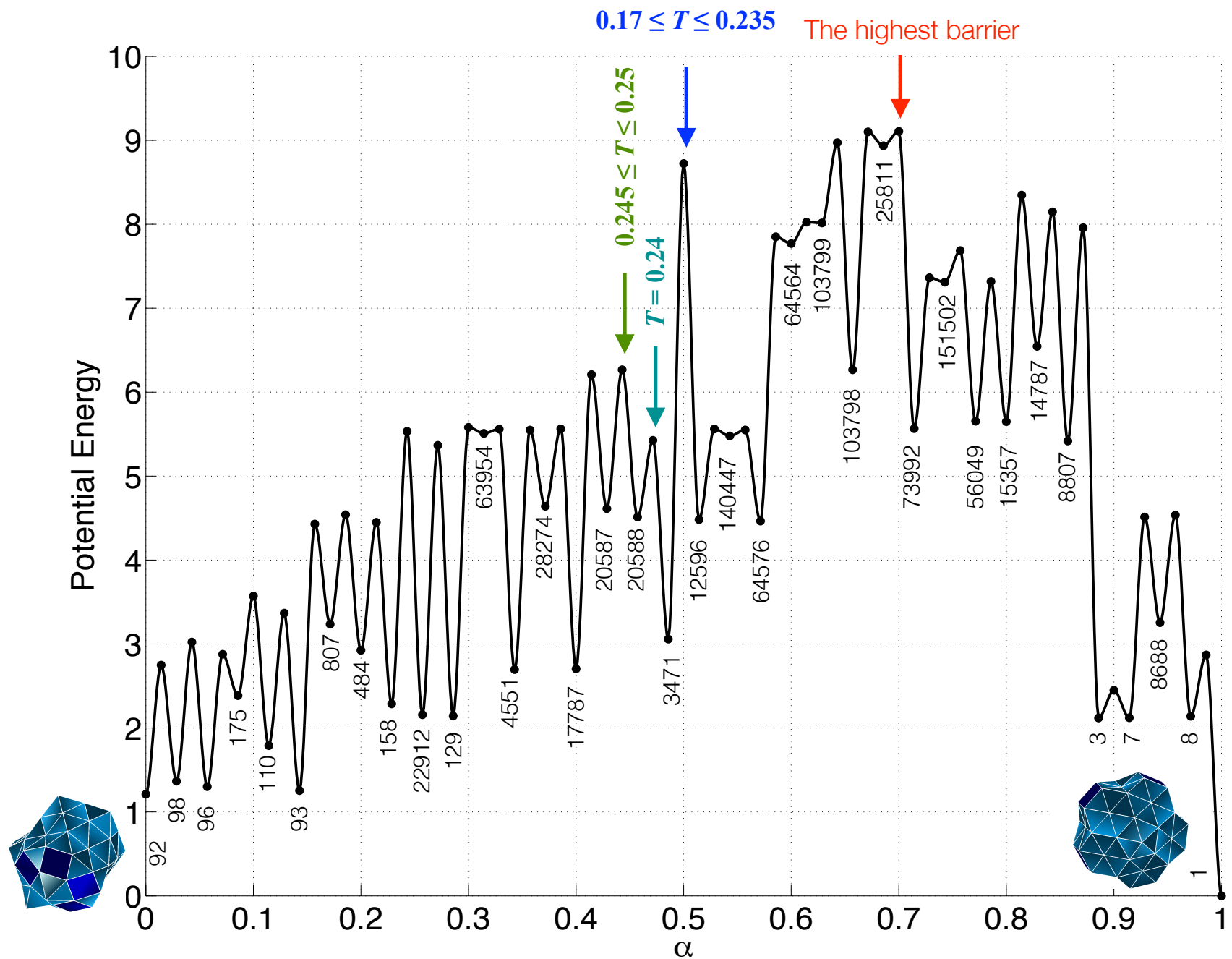
# Eigenvalue $\lambda_{4395}$ of $LJ_{75}$



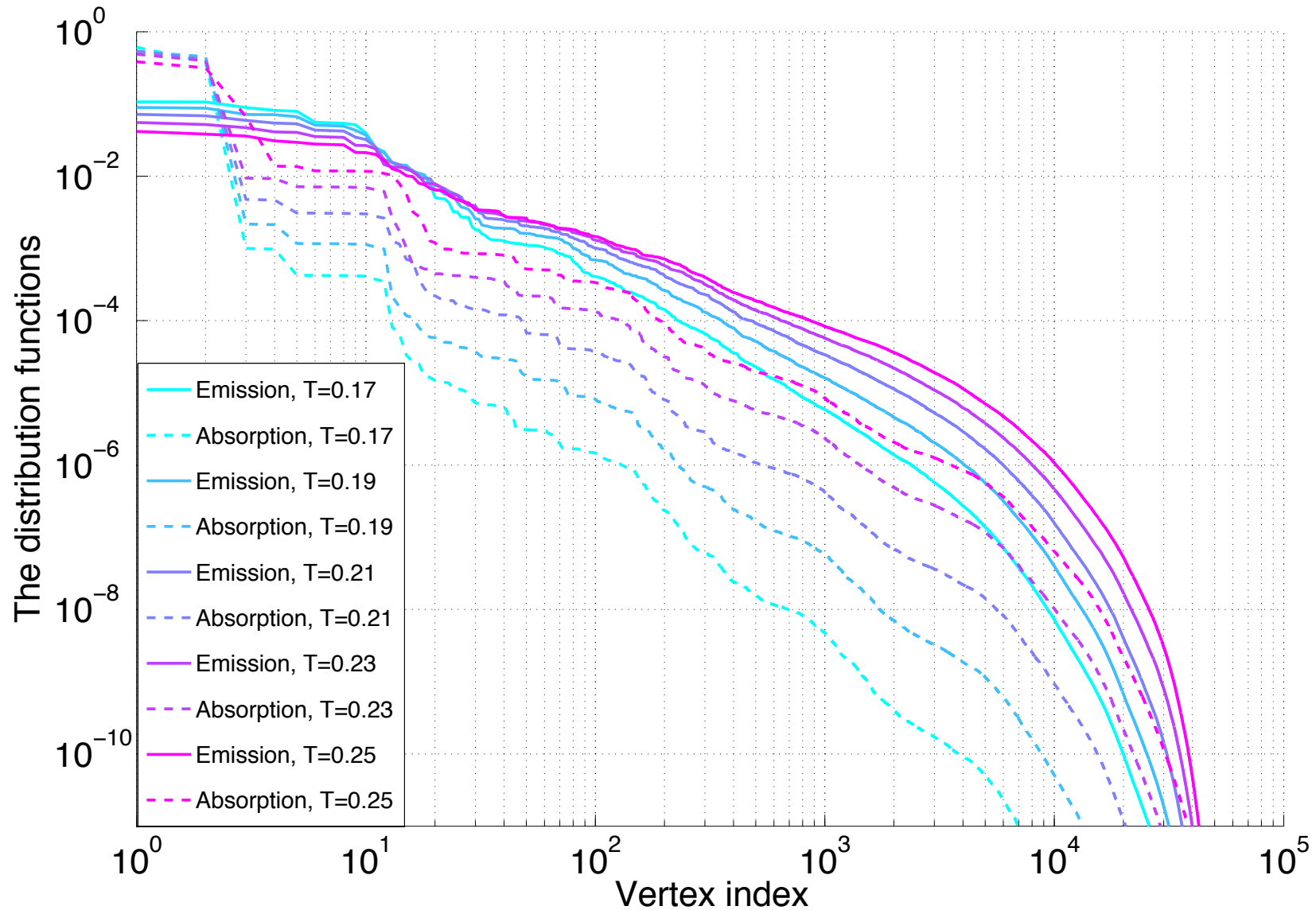
# Eigencurrent distribution in the emission-absorption cut



# The emission-absorption cut: location

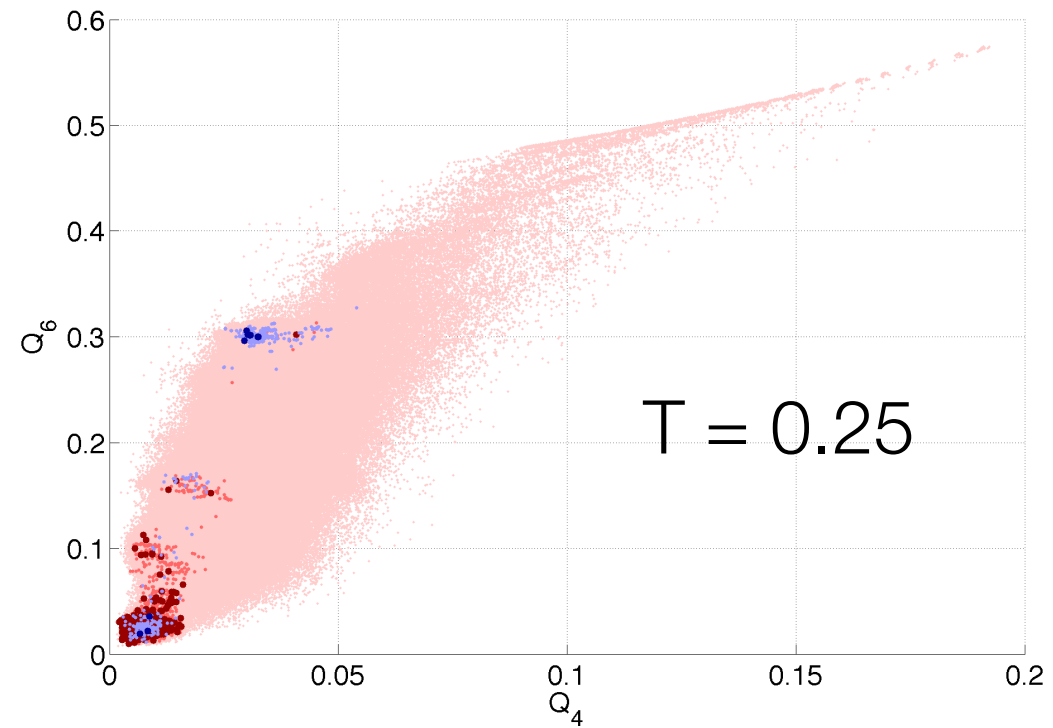
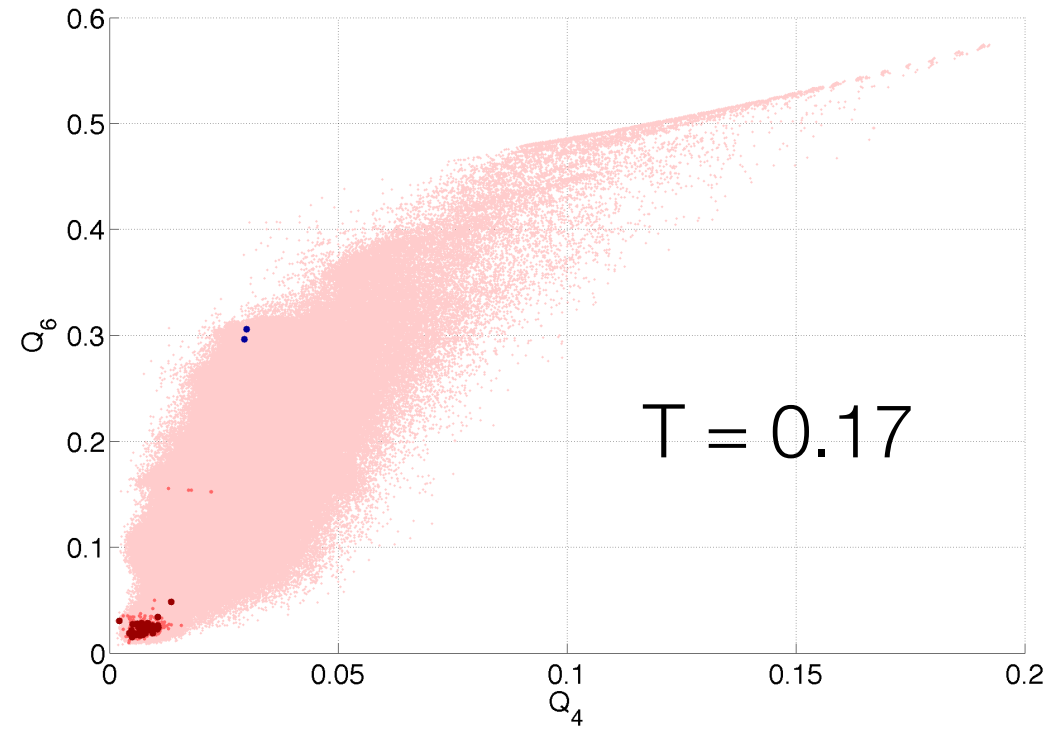


# Emission-absorption distribution



# Emission - absorption

- States emitting 99% of the eigencurrent
- States emitting 90% of the eigencurrent
- States absorbing 99% of the eigencurrent
- States absorbing 90% of the eigencurrent



# Highlights

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- Formulas for asymptotic eigenvalues: exponents and prefactors
- Nested property of optimal W-graphs for stochastic networks with pairwise transition rates of the form  $L_{ij} = k_{ij} \exp(-U_{ij}/T)$
- A single-sweep algorithm for computing asymptotic eigenvalues and eigenvectors
- Application to Lennard-Jones-75

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# Acknowledgements

- David Wales (Cambridge University, UK)